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How can we implement ontologies? Ontology languages

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Main References



Gómez-Pérez, A.; Fernández-López, M.; Corcho, O. *Ontological Engineering*. Springer Verlag, 2003



Baader F, McGuinness D, Nardi D, Patel-Schneider P (2003)
The Description Logic Handbook: Theory, implementation and applications.
Cambridge University Press, Cambridge, United Kingdom



<http://knowledgeweb.semanticweb.org>



Research deliverables
Industry deliverables



Dean M, Schreiber G (2004) *OWL Web Ontology Language Reference*. W3C Recommendation. <http://www.w3.org/TR/owl-ref/>
Brickley D, Guha RV (2004) *RDF Vocabulary Description Language 1.0: RDF Schema*. W3C Recommendation.
<http://www.w3.org/TR/PR-rdf-schema>
Lassila O, Swick R (1999) *Resource Description Framework (RDF) Model and Syntax Specification*. W3C Recommendation.
<http://www.w3.org/TR/REC-rdf-syntax/>



Acknowledgements

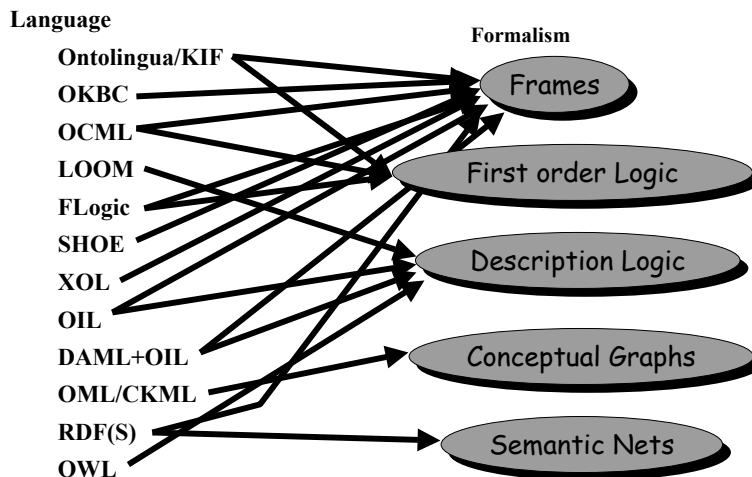
- **Asunción Gómez-Pérez and Mariano Fernández-López**
 - Most of the slides have been done jointly with them
- **Sean Bechhoffer (University of Manchester)**
 - Tableaux reasoning
 - Examples about reasoning
- **CO-ODE people (University of Manchester)**
 - <http://www.co-ode.org/>
 - Some RDF, RDFS and OWL descriptions
 - Use of reasoners
 - Protégé-OWL tutorial



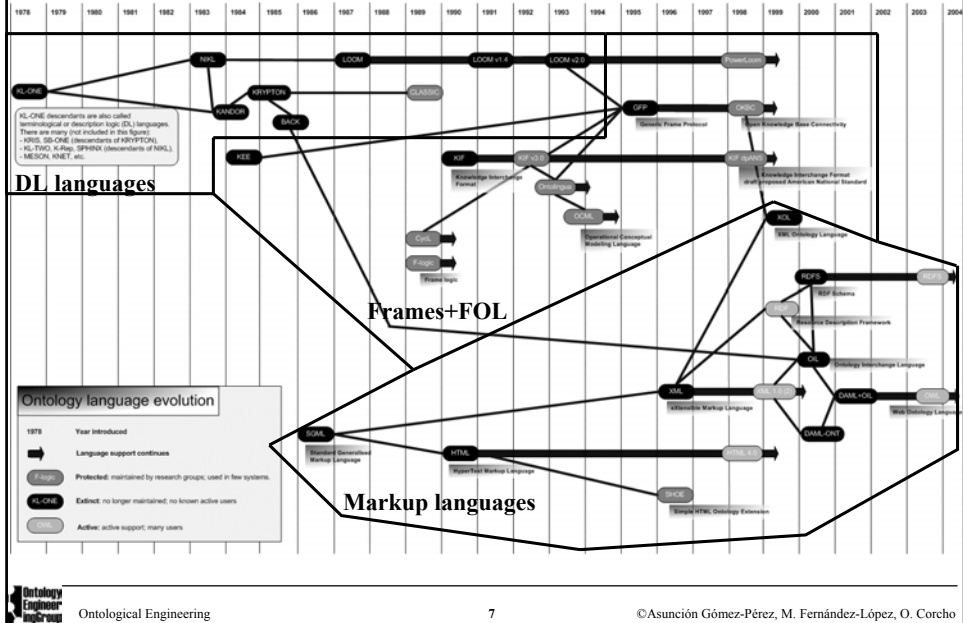
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KR Formalisms



Ontology language evolution



Ontology Languages (I)

Traditional ontology languages

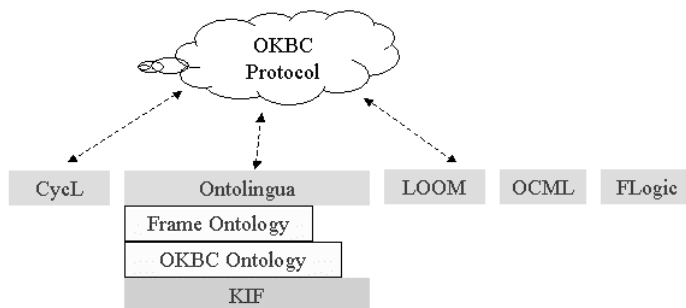
Ontolingua/KIF

OKBC

OCML

LOOM

FLogic



Ontology Languages (II)

Ontology markup languages

Standards & Recommendations of W3C

XML

RDF(S)

Ontology specification languages

SHOE

XOL

OIL

DAML+OIL

OWL

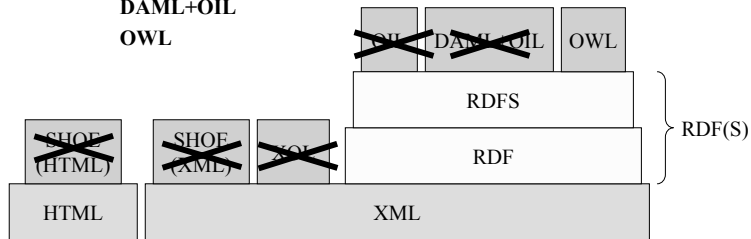
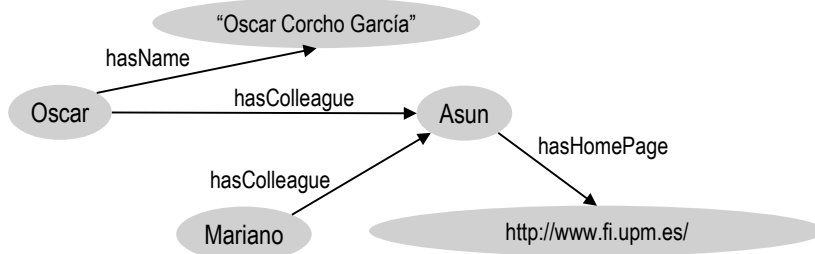


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RDF: Resource Description Framework

- **W3C recommendation** (<http://www.w3.org/RDF>)
- **RDF is graphical formalism** (+ XML syntax + semantics)
 - for representing metadata
 - for describing the semantics of information in a machine- accessible way
- **RDF is a basic ontology language**
 - Resources are described in terms of properties and property values using RDF statements.
 - Statements are represented as triples, consisting of a subject, predicate and object. [S, P, O]



RDF and URIs

- **RDF uses URIRefs (Unique Resource Identifiers References) to identify resources.**
 - A URIRef consists of a URI and an optional Fragment Identifier separated from the URI by the hash symbol #.
 - Examples
 - <http://www.co-ode.org/people#hasColleague>
 - `coode:hasColleague`
- **A set of URIRefs is known as a vocabulary**
 - E.g., the RDF Vocabulary
 - The set of URIRefs used in describing the RDF concepts: `rdf:Property`, `rdf:Resource`, `rdf:type`, etc.
 - The RDFS Vocabulary
 - The set of URIRefs used in describing the RDF Schema language: `rdfs:Class`, `rdfs:domain`, etc.
 - The ‘Pizza Ontology’ Vocabulary
 - `pz:hasTopping`, `pz:Pizza`, `pz:VegetarianPizza`, etc.

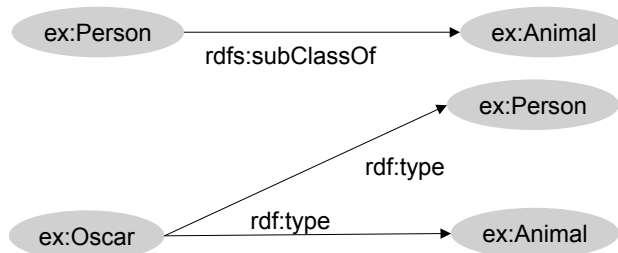
RDF Serialisation

```
<?xml version="1.0"?>
<rdf:RDF xmlns:rdf="http://www.w3.org/1999/02/22-rdf-syntax-ns#"
  xmlns:rdfs="http://www.w3.org/2000/01/rdf-schema#"
  xmlns:coode="http://www.co-ode.org/people#"
  xml:base="http://www.co-ode.org/people">
  <rdf:Description rdf:ID="mh">
    <coode:hasHomepage rdf:resource="http://www.cs.man.ac.uk/~horridgm"/>
    <coode:hasName>Matthew Horridge</coode:hasName>
  </rdf:Description>
  <rdf:Description rdf:ID="nd">
    <coode:hasName>Nick Drummond</coode:hasName>
    <coode:hasColleague rdf:resource="#mh"/>
  </rdf:Description>
</rdf:RDF>
```

```
<?xml version="1.0"?>
<rdf:RDF
  xmlns:coode="http://www.co-ode.org/people#"
  xmlns:rdf="http://www.w3.org/1999/02/22-rdf-syntax-ns#"
  xml:base="file:/Users/matthewhorridge/Desktop/Test.rdf">
  <rdf:Description rdf:about="http://www.co-ode.org/people#nd">
    <coode:hasName>Nick Drummond</coode:hasName>
    <coode:hasColleague>
      <rdf:Description rdf:about="http://www.co-ode.org/people#mh">
        <coode:hasName>Matthew Horridge</coode:hasName>
        <coode:hasHomepage rdf:resource="http://www.cs.man.ac.uk/~horridgm"/>
      </rdf:Description>
    </coode:hasColleague>
  </rdf:Description>
</rdf:RDF>
```

RDFS: RDF Schema

- W3C Recommendation
- RDF Schema extends RDF to enable talking about classes of resources, and the properties to be used with them.
 - Class definition: `rdfs:Class`, `rdfs:subClassOf`
 - Property definition: `rdfs:subPropertyOf`, `rdfs:range`, `rdfs:domain`
 - Other primitives: `rdfs:comment`, `rdfs:label`, `rdfs:seeAlso`, `rdfs:isDefinedBy`
- RDF Schema provides the means to describe application specific RDF vocabularies.
- RDFS vocabulary adds constraints on models, e.g.:
 - $\forall x,y,z$ `type(x,y)` and `subClassOf(y,z)` \rightarrow `type(x,z)`



RDF(S) limitations

- **RDFS too weak to describe resources in sufficient detail**
 - No localised range and domain constraints
 - Can't say that the range of hasChild is person when applied to persons and elephant when applied to elephants
 - No existence/cardinality constraints
 - Can't say that all *instances* of person have a mother that is also a person, or that persons have exactly 2 parents
 - No transitive, inverse or symmetrical properties
 - Can't say that isPartOf is a transitive property, that hasPart is the inverse of isPartOf or that touches is symmetrical
 - ...
- **Difficult to provide reasoning support**
 - No “native” reasoners for non-standard semantics
 - May be possible to reason via FO axiomatisation

Exercise



•Objective

- Understand the features of RDF(S) for implementing ontologies, including its limitations
- Get used to the graph and XML syntax of RDF(S)

•Tasks

- Take ontologies described in the previous day and create their graphs
 - First only include the vocabulary from the domain
 - Then include references to the RDF and RDFS vocabularies
- Serialise part of the graph in the XML syntax

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OWL

Web Ontology Language

Built on top of RDF(S) and renaming DAML+OIL primitives

3 layers:

- OWL Lite: a small subset, easier for frame-based tools to transition to, easier reasoning
- OWL DL: description logic, decidable reasoning
- OWL Full: RDF extension, allows metaclasses

Several syntaxes:

- Abstract syntax: easier to read and write manually, closely corresponds to DL
- RDF/XML: OWL can be parsed as an RDF document, more verbose



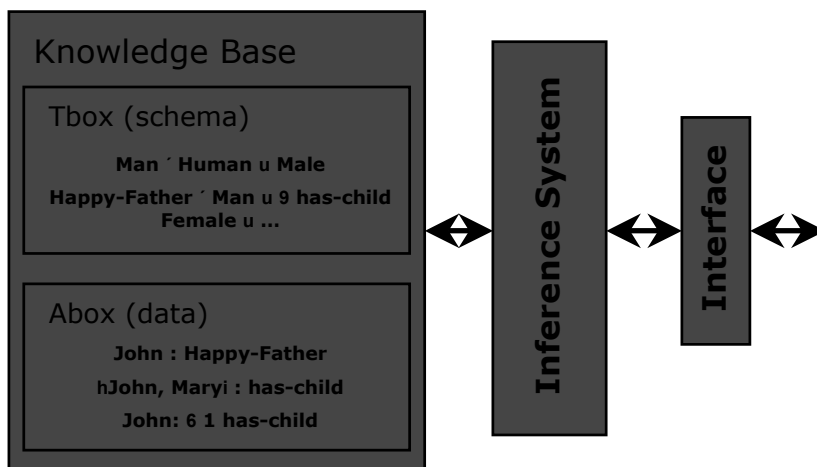
Dean M, Schreiber G. The OWL Web Ontology Language. W3C Recommendation. February 2004.



OWL and Description Logic

- **A family of logic based Knowledge Representation formalisms**
 - Descendants of semantic networks and KL-ONE
 - Describe domain in terms of concepts (classes), roles (relationships) and individuals
 - Specific languages characterised by the constructors and axioms used to assert knowledge about classes, roles and individuals.
 - Example: ALC (the least expressive language in DL that is propositionally closed)
 - Constructors: boolean (and, or, not)
 - Role restrictions
- **Distinguished by:**
 - Formal semantics (typically model theoretic)
 - Decidable fragments of FOL
 - Closely related to Propositional Modal & Dynamic Logics
 - Provision of inference services
 - Sound and complete decision procedures for key problems
 - Implemented systems (highly optimised)

DL Architecture



DL constructors

Construct	Syntax	Language		
Concept	A	FL ₀	AL	
Role name	R			
Intersection	$C \cap D$	S ¹⁴		
Value restriction	$\forall R.C$			
Limited existential quantification	$\exists R$			
Top or Universal	\top			
Bottom	\perp			
Atomic negation	$\neg A$			
Negation ¹⁵	$\neg C$			C
Union	$C \cup D$			U
Existential restriction	$\exists R.C$		E	
Number restrictions	$(\geq n R) (\leq n R)$		N	
Nominals	$\{a_1 \dots a_n\}$	O		
Role hierarchy	$R \subseteq S$	H		
Inverse role	R^{-}	I		
Qualified number restriction	$(\geq n R.C) (\leq n R.C)$	Q		

OWL is SHOIN(D+)

Other:

Concrete datatypes: hasAge (<21)

Transitive roles: hasChild* (descendant)

Role composition: hasParent o hasBrother (uncle)

¹² Names previously used for Description Logics were: terminological knowledge representation languages, concept languages, term subsumption languages, and KL-ONE-based knowledge representation languages.

¹³ In this table, we use A to refer to atomic concepts (concepts that are the basis for building other concepts), C and D to any concept definition, R to atomic roles and S to role definitions. FL is used for structural DL languages and AL for attributive languages (Baader et al., 2003).

¹⁴ S is the name used for the language ALC_{S^*} , which is composed of ALC plus transitive roles.

¹⁵ ALC and ALCUE are equivalent languages, since union (U) and existential restriction (E) can be represented using negation (C).



OWL as a Description Logic language. Class constructors

Constructor	DL Syntax	Example	Modal Syntax
intersectionOf	$C_1 \sqcap \dots \sqcap C_n$	Human \sqcap Male	$C_1 \wedge \dots \wedge C_n$
unionOf	$C_1 \sqcup \dots \sqcup C_n$	Doctor \sqcup Lawyer	$C_1 \vee \dots \vee C_n$
complementOf	$\neg C$	\neg Male	$\neg C$
oneOf	$\{x_1\} \sqcup \dots \sqcup \{x_n\}$	{john} \sqcup {mary}	$x_1 \vee \dots \vee x_n$
allValuesFrom	$\forall P.C$	\forall hasChild.Doctor	$[P]C$
someValuesFrom	$\exists P.C$	\exists hasChild.Lawyer	$\langle P \rangle C$
maxCardinality	$\leq n P$	≤ 1 hasChild	$[P]_{n+1}$
minCardinality	$\geq n P$	≥ 2 hasChild	$\langle P \rangle_n$

- **XML Schema datatypes are treated as classes**
 - \forall hasAge.nonNegativeInteger
- **Nesting of constructors can be arbitrarily complex**
 - $\text{Person} \wedge \forall \text{hasChild} . (\text{Doctor} \vee \exists \text{hasChild} . \text{Doctor})$
- **Lots of redundancy, e.g., use negations to transform and to or and exists to forall**



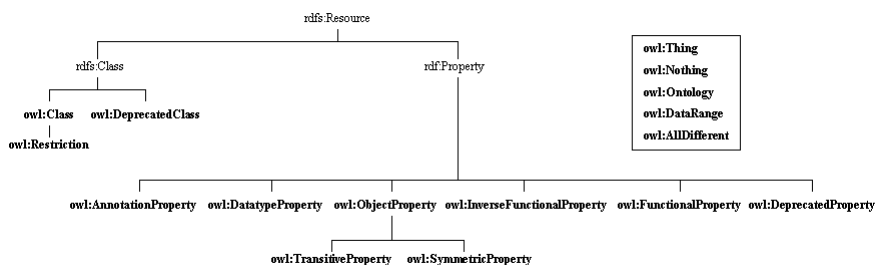
OWL Axioms

Axiom	DL Syntax	Example
subClassOf	$C_1 \sqsubseteq C_2$	Human \sqsubseteq Animal \sqcap Biped
equivalentClass	$C_1 \equiv C_2$	Man \equiv Human \sqcap Male
disjointWith	$C_1 \sqsubseteq \neg C_2$	Male \sqsubseteq \neg Female
sameIndividualAs	$\{x_1\} \equiv \{x_2\}$	{President_Bush} \equiv {G_W_Bush}
differentFrom	$\{x_1\} \sqsubseteq \neg\{x_2\}$	{john} \sqsubseteq \neg {peter}
subPropertyOf	$P_1 \sqsubseteq P_2$	hasDaughter \sqsubseteq hasChild
equivalentProperty	$P_1 \equiv P_2$	cost \equiv price
inverseOf	$P_1 \equiv P_2^-$	hasChild \equiv hasParent $^-$
transitiveProperty	$P^+ \sqsubseteq P$	ancestor $^+$ \sqsubseteq ancestor
functionalProperty	$\top \sqsubseteq \leq 1P$	$\top \sqsubseteq \leq 1$ hasMother
inverseFunctionalProperty	$\top \sqsubseteq \leq 1P^-$	$\top \sqsubseteq \leq 1$ hasSSN $^-$

- **Axioms (mostly) reducible to inclusion (v)**

- $C \equiv D$ iff both $C \sqsubseteq D$ and $D \sqsubseteq C$
- C disjoint D iff $C \sqsubseteq \neg D$

Class taxonomy of the OWL KR ontology



Property list of the OWL KR ontology

Property name	domain	range
owl:intersectionOf	owl:Class	rdf:List
owl:unionOf	owl:Class	rdf:List
owl:complementOf	owl:Class	owl:Class
owl:oneOf	owl:Class	rdf:List
owl:onProperty	owl:Restriction	rdf:Property
owl:allValuesFrom	owl:Restriction	rdfs:Class
owl:hasValue	owl:Restriction	<i>not specified</i>
owl:someValuesFrom	owl:Restriction	rdfs:Class
owl:minCardinality	owl:Restriction	xsd:nonNegativeInteger OWL Lite: {0,1} OWL DL/Full: {0,...,N}
owl:maxCardinality	owl:Restriction	xsd:nonNegativeInteger OWL Lite: {0,1} OWL DL/Full: {0,...,N}
owl:cardinality	owl:Restriction	xsd:nonNegativeInteger OWL Lite: {0,1} OWL DL/Full: {0,...,N}
owl:inverseOf	owl:ObjectProperty	owl:ObjectProperty
owl:sameAs	owl:Thing	owl:Thing
owl:equivalentClass	owl:Class	owl:Class
owl:equivalentProperty	rdf:Property	rdf:Property
owl:sameIndividualAs	owl:Thing	owl:Thing
owl:differentFrom	owl:Thing	owl:Thing
owl:disjointWith	owl:Class	owl:Class
owl:distinctMembers	owl:AllDifferent	rdf:List
owl:versionInfo	<i>not specified</i>	<i>not specified</i>
owl:priorVersion	owl:Ontology	owl:Ontology
owl:incompatibleWith	owl:Ontology	owl:Ontology
owl:backwardCompatibleWith	owl:Ontology	owl:Ontology
owl:imports	owl:Ontology	owl:Ontology



<p>OWL DL Class expressions allowed in: <i>rdfs:domain, rdfs:range, rdfs:subClassOf, owl:intersectionOf, owl:equivalentClass, owl:hasValueFrom, owl:someValuesFrom</i> Values are not restricted (0..N) in: <i>owl:minCardinality, owl:maxCardinality, owl:cardinality</i></p> <p>owl:DataRange, <i>rdfs:List, rdfs:first, rdfs:rest, rdfs:tail</i></p> <p>owl:hasValue (<i>owl:hasValue</i>) owl:oneOf (<i>owl:oneOf</i>) owl:unionOf (<i>owl:unionOf</i>), owl:complementOf (<i>owl:complementOf</i>) owl:disjointWith (<i>owl:disjointWith</i>)</p> <p>OWL Lite owl:Ontology (<i>owl:Ontology</i>), owl:versionInfo (<i>owl:versionInfo</i>), owl:imports (<i>owl:imports</i>), owl:backwardCompatibleWith, owl:incompatibleWith, owl:priorVersion, owl:DeprecatedClass, owl:DeprecatedProperty</p> <p>owl:Class (<i>owl:Class</i>), owl:Restriction (<i>owl:Restriction</i>), owl:onProperty (<i>owl:onProperty</i>), owl:allValuesFrom (<i>owl:allValuesFrom</i>) (only with class identifiers and named datatypes), owl:someValuesFrom (<i>owl:someValuesFrom</i>) (only with class identifiers and named datatypes), owl:minCardinality (<i>owl:minCardinality</i>, restricted to (0,1)), owl:maxCardinality (<i>owl:maxCardinality</i>, restricted to (0,1)), owl:cardinality (<i>owl:cardinality</i>, restricted to (0,1))</p> <p>owl:intersectionOf (only with class identifiers and property restrictions)</p> <p>owl:ObjectProperty (<i>owl:ObjectProperty</i>), owl:DatatypeProperty (<i>owl:DatatypeProperty</i>), owl:TransitiveProperty (<i>owl:TransitiveProperty</i>), owl:SymmetricProperty, owl:FunctionalProperty (<i>owl:FunctionalProperty</i>), owl:InverseFunctionalProperty (<i>owl:InverseFunctionalProperty</i>), owl:AnnotationProperty</p> <p>owl:Thing (<i>owl:Thing</i>) owl:Nothing (<i>owl:Nothing</i>)</p> <p>owl:inverseOf (<i>owl:inverseOf</i>), owl:equivalentClass (<i>owl:equivalentClass</i>) (only with class identifiers and property restrictions), owl:equivalentProperty (<i>owl:equivalentProperty</i>), owl:sameAs (<i>owl:sameAs</i>), owl:sameIndividualAs, owl:differentFrom (<i>owl:differentFrom</i>), owl:AllDifferent, owl:distinctMembers</p> <p>RDFS rdf:Property rdfs:subPropertyOf rdfs:domain rdfs:range (only with class identifiers and named datatypes) rdfs:comment, rdfs:label, rdfs:seeAlso, rdfs:isDefinedBy rdfs:subClassOf (only with class identifiers and property restrictions)</p>



OWL DL
Class expressions allowed in: `rdfs:domain`, `rdfs:range`, `rdfs:subClassOf`, `owl:intersectionOf`, `owl:equivalentClass`, `owl:allValuesFrom`, `owl:someValuesFrom`, `owl:mainCardinality`, `owl:maxCardinality`, `owl:cardinality`

Values are not restricted (0..1) in:
`owl:DataRange`, `rdfs:List`, `rdfs:first`, `rdfs:rest`, `rdfs:tail`

`owl:hasValue` (*daml:hasValue*)
`owl:oneOf` (*daml:oneOf*)
`owl:unionOf` (*daml:unionOf*), `owl:complementOf` (*daml:complementOf*)
`owl:disjointWith` (*daml:disjointWith*)

OWL Lite
`owl:Ontology` (*daml:Ontology*),
`owl:versionInfo` (*daml:versionInfo*),
`owl:imports` (*daml:imports*),
`owl:backwardCompatibleWith`,
`owl:incompatibleWith`, `owl:priorVersion`,
`owl:DeprecatedClass`,
`owl:DeprecatedProperty`

`owl:Class` (*daml:Class*),
`owl:Restriction` (*daml:Restriction*),
`owl:Property` (*daml:Property*),
`owl:allValuesFrom` (*daml:toClass*) (only with class identifiers and named datatypes),
`owl:someValuesFrom` (*daml:hasClass*) (only with class identifiers and named datatypes),
`owl:mainCardinality` (*daml:minCardinality*, restricted to {0,1}),
`owl:maxCardinality` (*daml:maxCardinality*, restricted to {0,1}),
`owl:cardinality` (*daml:cardinality*, restricted to {0,1})

`owl:intersectionOf` (only with class identifiers and property restrictions)

`owl:ObjectProperty` (*daml:ObjectProperty*),
`owl:DatatypeProperty` (*daml:DatatypeProperty*),
`owl:TransitiveProperty` (*daml:TransitiveProperty*),
`owl:SymmetricProperty`,
`owl:FunctionalProperty` (*daml:UniqueProperty*),
`owl:InverseFunctionalProperty` (*daml:UnambiguousProperty*),
`owl:AnnotationProperty`

`owl:Thing` (*daml:Thing*)
`owl:Nothing` (*daml:Nothing*)

`owl:inverseOf` (*daml:inverseOf*),
`owl:equivalentClass` (*daml:sameClassAs*) (only with class identifiers and property restrictions),
`owl:equivalentProperty` (*daml:samePropertyAs*),
`owl:sameAs` (*daml:equivalentTo*),
`owl:sameIndividualAs`,
`owl:differentFrom` (*daml:differentIndividualFrom*),
`owl:AllDifferent`, `owl:distinctMembers`

RDF(S)
`rdf:Property`
`rdfs:subPropertyOf`
`rdfs:domain`
`rdfs:range` (only with class identifiers and named datatypes)
`rdfs:comment`, `rdfs:label`, `rdfs:seeAlso`, `rdfs:isDefinedBy`
`rdfs:subClassOf` (only with class identifiers and property restrictions)

RDF(S)
`rdf:Property`
`rdfs:subPropertyOf`
`rdfs:domain`
`rdfs:range` (only with class identifiers and named datatypes)
`rdfs:comment`, `rdfs:label`, `rdfs:seeAlso`, `rdfs:isDefinedBy`
`rdfs:subClassOf` (only with class identifiers and property restrictions)

R

$R \subseteq S$

$C \subseteq D$

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OWL DL
Class expressions allowed in: `rdfs:domain`, `rdfs:range`, `rdfs:subClassOf`, `owl:intersectionOf`, `owl:equivalentClass`, `owl:allValuesFrom`, `owl:someValuesFrom`, `owl:mainCardinality`, `owl:maxCardinality`, `owl:cardinality`

Values are not restricted (0..1) in:
`owl:DataRange`, `rdfs:List`, `rdfs:first`, `rdfs:rest`, `rdfs:tail`

`owl:hasValue` (*daml:hasValue*)
`owl:oneOf` (*daml:oneOf*)
`owl:unionOf` (*daml:unionOf*), `owl:complementOf` (*daml:complementOf*)
`owl:disjointWith` (*daml:disjointWith*)

OWL Lite
`owl:Ontology` (*daml:Ontology*),
`owl:versionInfo` (*daml:versionInfo*),
`owl:imports` (*daml:imports*),
`owl:backwardCompatibleWith`,
`owl:incompatibleWith`, `owl:priorVersion`,
`owl:DeprecatedClass`,
`owl:DeprecatedProperty`

`owl:Class` (*daml:Class*),
`owl:Restriction` (*daml:Restriction*),
`owl:Property` (*daml:Property*),
`owl:allValuesFrom` (*daml:toClass*) (only with class identifiers and named datatypes),
`owl:someValuesFrom` (*daml:hasClass*) (only with class identifiers and named datatypes),
`owl:mainCardinality` (*daml:minCardinality*, restricted to {0,1}),
`owl:maxCardinality` (*daml:maxCardinality*, restricted to {0,1}),
`owl:cardinality` (*daml:cardinality*, restricted to {0,1})

`owl:intersectionOf` (only with class identifiers and property restrictions)

`owl:ObjectProperty` (*daml:ObjectProperty*),
`owl:DatatypeProperty` (*daml:DatatypeProperty*),
`owl:TransitiveProperty` (*daml:TransitiveProperty*),
`owl:SymmetricProperty`,
`owl:FunctionalProperty` (*daml:UniqueProperty*),
`owl:InverseFunctionalProperty` (*daml:UnambiguousProperty*),
`owl:AnnotationProperty`

`owl:Thing` (*daml:Thing*)
`owl:Nothing` (*daml:Nothing*)

`owl:inverseOf` (*daml:inverseOf*),
`owl:equivalentClass` (*daml:sameClassAs*) (only with class identifiers and property restrictions),
`owl:equivalentProperty` (*daml:samePropertyAs*),
`owl:sameAs` (*daml:equivalentTo*),
`owl:sameIndividualAs`,
`owl:differentFrom` (*daml:differentIndividualFrom*),
`owl:AllDifferent`, `owl:distinctMembers`

RDF(S)
`rdf:Property`
`rdfs:subPropertyOf`
`rdfs:domain`
`rdfs:range` (only with class identifiers and named datatypes)
`rdfs:comment`, `rdfs:label`, `rdfs:seeAlso`, `rdfs:isDefinedBy`
`rdfs:subClassOf` (only with class identifiers and property restrictions)

OWL Lite
`owl:Ontology` (*daml:Ontology*),
`owl:versionInfo` (*daml:versionInfo*),
`owl:imports` (*daml:imports*),
`owl:backwardCompatibleWith`,
`owl:incompatibleWith`, `owl:priorVersion`,
`owl:DeprecatedClass`,
`owl:DeprecatedProperty`

$\forall R.C$

$\exists R.C$

$\leq nR$

$\geq nR$

$= nR$

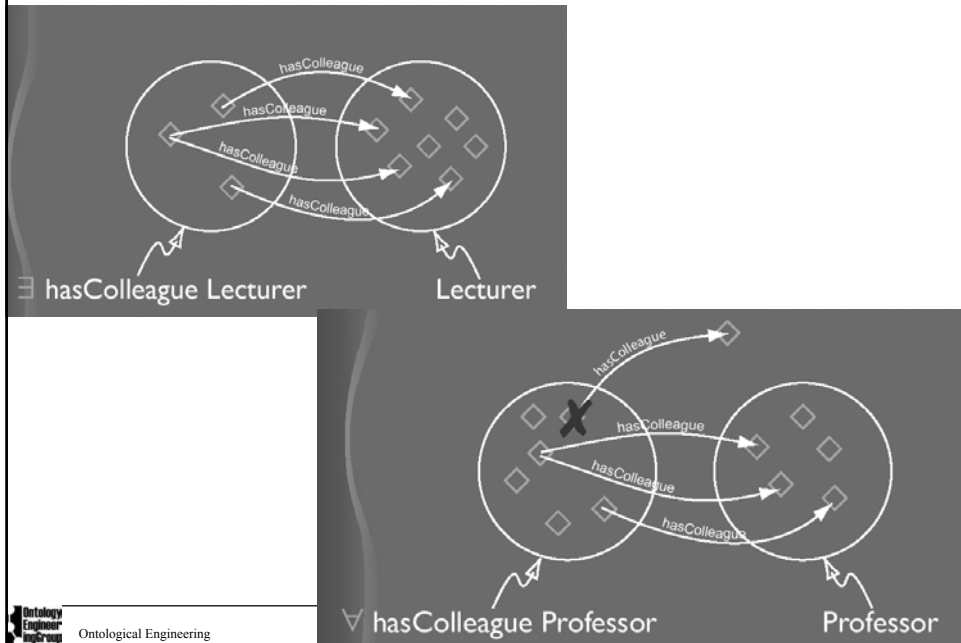
$C \cap D$

`owl:Class` (*daml:Class*),
`owl:Restriction` (*daml:Restriction*)
`owl:onProperty` (*daml:onProperty*),
`owl:allValuesFrom` (*daml:toClass*) (only with class identifiers and named datatypes),
`owl:someValuesFrom` (*daml:hasClass*) (only with class identifiers and named datatypes),
`owl:minCardinality` (*daml:minCardinality*, restricted to {0,1}),
`owl:maxCardinality` (*daml:maxCardinality*, restricted to {0,1}),
`owl:cardinality` (*daml:cardinality*, restricted to {0,1})

`owl:intersectionOf` (only with class identifiers and property restrictions)

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Existential and Universal Restrictions



Ontological Engineering

<p>OWL DL Class expressions allowed in: <i>rdfs:domain</i>, <i>rdfs:range</i>, <i>rdfs:subClassOf</i> Values are not restricted (0..N) in: <i>owl:intersectionOf</i>, <i>owl:equivalentClass</i>, <i>owl:allValuesFrom</i>, <i>owl:someValuesFrom</i> <i>owl>DataRange</i>, <i>rdfs:List</i>, <i>rdfs:first</i>, <i>rdfs:rest</i>, <i>rdfs:tail</i> <i>owl:hasValue</i> (<i>daml:hasValue</i>) <i>owl:oneOf</i> (<i>daml:oneOf</i>) <i>owl:isac</i> (<i>daml:isac</i>), <i>owl:complementOf</i> (<i>daml:complementOf</i>) <i>owl:disjointWith</i> (<i>daml:disjointWith</i>)</p>	<p><i>owl:ObjectProperty</i> (<i>daml:ObjectProperty</i>), <i>owl:DatatypeProperty</i> (<i>daml:DatatypeProperty</i>), <i>owl:TransitiveProperty</i> (<i>daml:TransitiveProperty</i>), $(+)$ R <i>owl:SymmetricProperty</i>, <i>owl:FunctionalProperty</i> (<i>daml:UniqueProperty</i>), <i>owl:InverseFunctionalProperty</i> (<i>daml:UnambiguousProperty</i>), <i>owl:AnnotationProperty</i></p>
<p>OWL Lite <i>owl:Ontology</i> (<i>daml:Ontology</i>), <i>owl:versionInfo</i> (<i>daml:versionInfo</i>), <i>owl:imports</i> (<i>daml:imports</i>), <i>owl:backwardsCompatibleWith</i>, <i>owl:isCompatibleWith</i>, <i>owl:priorVersion</i>, <i>owl:deprecatedClass</i>, <i>owl:deprecatedProperty</i> <i>owl:Class</i> (<i>daml:Class</i>), <i>owl:Restriction</i> (<i>daml:Restriction</i>), <i>owl:ObjectProperty</i> (<i>daml:ObjectProperty</i>), <i>owl:DatatypeProperty</i> (<i>daml:DatatypeProperty</i>), <i>owl:TransitiveProperty</i> (<i>daml:TransitiveProperty</i>), <i>owl:SymmetricProperty</i>, <i>owl:FunctionalProperty</i> (<i>daml:UniqueProperty</i>), <i>owl:InverseFunctionalProperty</i> (<i>daml:UnambiguousProperty</i>), <i>owl:AnnotationProperty</i></p>	<p><i>owl:Thing</i> (<i>daml:THING</i>) T <i>owl:Nothing</i> (<i>daml:Nothing</i>) \perp</p>
<p><i>owl:intersectionOf</i> (only with class identifiers and property restrictions) <i>owl:ObjectProperty</i> (<i>daml:ObjectProperty</i>), <i>owl:DatatypeProperty</i> (<i>daml:DatatypeProperty</i>), <i>owl:TransitiveProperty</i> (<i>daml:TransitiveProperty</i>), <i>owl:SymmetricProperty</i>, <i>owl:FunctionalProperty</i> (<i>daml:UniqueProperty</i>), <i>owl:InverseFunctionalProperty</i> (<i>daml:UnambiguousProperty</i>), <i>owl:AnnotationProperty</i></p>	<p>R^{-1} $C \equiv D$ $R \equiv S$</p>
<p><i>owl:Thing</i> (<i>daml:Thing</i>) <i>owl:Nothing</i> (<i>daml:Nothing</i>) <i>owl:inverseOf</i> (<i>daml:inverseOf</i>), <i>owl:equivalentClass</i> (<i>daml:sameClassAs</i>) (only with class identifiers and property restrictions), <i>owl:equivalentProperty</i> (<i>daml:samePropertyAs</i>), <i>owl:sameAs</i> (<i>daml:equivalentTo</i>), <i>owl:sameIndividualAs</i>, <i>owl:differentFrom</i> (<i>daml:differentIndividualFrom</i>), <i>owl:AllDifferent</i>, <i>owl:distinctMembers</i></p>	<p><i>owl:inverseOf</i> (<i>daml:inverseOf</i>), <i>owl:equivalentClass</i> (<i>daml:sameClassAs</i>) (only with class identifiers and property restrictions), <i>owl:equivalentProperty</i> (<i>daml:samePropertyAs</i>), <i>owl:sameAs</i> (<i>daml:equivalentTo</i>), <i>owl:sameIndividualAs</i>, <i>owl:differentFrom</i> (<i>daml:differentIndividualFrom</i>), <i>owl:AllDifferent</i>, <i>owl:distinctMembers</i> } $A \text{ box}$</p>
<p>RDFS(S) <i>rdfs:Property</i> <i>rdfs:subPropertyOf</i> <i>rdfs:domain</i> <i>rdfs:range</i> (only with class identifiers and named datatypes) <i>rdfs:comment</i>, <i>rdfs:label</i>, <i>rdfs:seeAlso</i>, <i>rdfs:isDefinedBy</i> <i>rdfs:subClassOf</i> (only with class identifiers and property restrictions)</p>	<p>©Asunción Gómez-Pérez, M. Fernández-López, O. Corcho</p>

<p>OWL DL</p> <p>Class expressions allowed in: rdfs:domain, rdfs:range, rdfs:subClassOf owl:intersectionOf, owl:equivalentClass, owl:allValuesFrom, owl:someValuesFrom owl:minCardinality, owl:maxCardinality, owl:cardinality</p> <p>Values are not restricted (0..N) in:</p> <p>owl:DataRange, rdf:List, rdf:first, rdf:rest, rdf:nil</p> <p>owl:hasValue (daml:hasValue) owl:oneOf (daml:oneOf) owl:unionOf (daml:unionOf), owl:complementOf (daml:complementOf) owl:disjointWith (daml:disjointWith)</p> <p>OWL Lite</p> <p>owl:Ontology (daml:Ontology), owl:versionInfo (daml:versionInfo), owl:imports (daml:imports), owl:backwardCompatibleWith, owl:incompatibleWith, owl:priorVersion, owl:DeprecatedClass, owl:DeprecatedProperty</p> <p>RDFS</p> <p>rdf:Property rdfs:subPropertyOf rdfs:domain rdfs:range (only with class identifiers and named datatypes) rdfs:comment, rdfs:label, rdfs:seeAlso, rdfs:isDefinedBy rdfs:subClassOf (only with class identifiers and property restrictions)</p>	<p>Class expressions allowed in: rdfs:domain, rdfs:range, rdfs:subClassOf owl:intersectionOf, owl:equivalentClass, owl:allValuesFrom, owl:someValuesFrom</p> <p>Values are not restricted (0..N) in: owl:minCardinality, owl:maxCardinality, owl:cardinality</p> <p>owl:DataRange, rdf:List, rdf:first, rdf:rest, rdf:nil</p> <p>owl:hasValue (daml:hasValue) → $\exists R. \{x\}$</p> <p>owl:oneOf (daml:oneOf) → $\{x\}$</p> <p>owl:unionOf (daml:unionOf), owl:complementOf (daml:complementOf) → $\neg C$</p> <p>owl:disjointWith (daml:disjointWith) → $C \cup D$</p> <p>owl:Nothing (daml:Nothing) → $C \cap D \subseteq \perp$</p>
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hasValue and oneOf

hasColleague \ni Matthew Person

HolidayDestinations

To specify an enumerated class, the individuals that are members of the class are listed inside curly brackets {...}

{Spain Germany France Italy}

OWL Example



Develop a sample ontology in the domain of people, pets, vehicles, and newspapers

- Practice with DL syntax, OWL abstract syntax and OWL RDF/XML syntax
- Understand the basic primitives of OWL Lite and OWL DL
- Understand the basic reasoning mechanisms of OWL DL (tomorrow)

Subsumption

Automatic classification: an ontology built collaboratively

Instance classification

Detecting redundancy

Consistency checking: unsatisfiable restrictions in a Tbox (are the classes coherent?)

Some basic DL modelling guidelines

- X must be Y, X is an Y that... $\rightarrow X \subseteq Y$
- X is exactly Y, X is the Y that... $\rightarrow X \equiv Y$
- X is not Y $\rightarrow X \subseteq \neg Y$
- X and Y are disjoint $\rightarrow X \cap Y \subseteq \perp$
- X is Y or Z $\rightarrow X \subseteq Y \cup Z$
- X is Y for which property P has only instances of Z as values $\rightarrow X \subseteq Y \cap (\forall P.Z)$
- X is Y for which property P has at least an instance of Z as a value $\rightarrow X \subseteq Y \cap (\exists P.Z)$
- X is Y for which property P has at most 2 values $\rightarrow X \subseteq Y \cap (\leq 2.P)$
- Individual X is a Y $\rightarrow X \in Y$

Chunk 1. Formalize in DL, and then in OWL DL

1. Concept definitions:

Grass and trees must be plants. Leaves are parts of a tree but there are other parts of a tree that are not leaves. A dog must eat bones, at least. A sheep is an animal that must only eat grass. A giraffe is an animal that must only eat leaves. A mad cow is a cow that eats brains that can be part of a sheep.

2. Restrictions:

Animals or part of animals are disjoint with plants or parts of plants.

3. Properties:

Eats is applied to animals. Its inverse is eaten_by.

4. Individuals:

Tom.

Flossie is a cow.

Rex is a dog and is a pet of Mick.

Fido is a dog.

Tibbs is a cat.

Chunk 2. Formalize in DL, and then in OWL DL

1. Concept definitions:

Bicycles, buses, cars, lorries, trucks and vans are vehicles. There are several types of companies: bus companies and haulage companies.

An elderly person must be adult. A kid is (exactly) a person who is young. A man is a person who is male and is adult. A woman is a person who is female and is adult. A grown up is a person who is an adult. And old lady is a person who is elderly and female. Old ladies must have some animal as pets and all their pets are cats.

2. Restrictions:

Youngs are not adults, and adults are not youngs.

3. Properties:

Has mother and has father are subproperties of has parent.

4. Individuals:

Kevin is a person.

Fred is a person who has a pet called Tibbs.

Joe is a person who has at most one pet. He has a pet called Fido.

Minnie is a female, elderly, who has a pet called Tom.

Chunk 3. Formalize in DL, and then in OWL DL

1. Concept definitions:

A magazine is a publication. Broadsheets and tabloids are newspapers. A quality broadsheet is a type of broadsheet. A red top is a type of tabloid. A newspaper is a publication that must be either a broadsheet or a tabloid.

White van mans must read only tabloids.

2. Restrictions:

Tabloids are not broadsheets, and broadsheets are not tabloids.

3. Properties:

The only things that can be read are publications.

4. Individuals:

Daily Mirror

The Guardian and The Times are broadsheets

The Sun is a tabloid

Chunk 4. Formalize in DL, and then in OWL DL

1. Concept definitions:

A pet is a pet of something. An animal must eat something. A vegetarian is an animal that does not eat animals nor parts of animals. Ducks, cats and tigers are animals.

An animal lover is a person who has at least three pets. A pet owner is a person who has animal pets. A cat liker is a person who likes cats. A cat owner is a person who has cat pets. A dog liker is a person who likes dogs. A dog owner is a person who has dog pets.

2. Restrictions:

Dogs are not cats, and cats are not dogs.

3. Properties:

Has pet is defined between persons and animals. Its inverse is is_pet_of.

4. Individuals:

Dewey, Huey, and Louie are ducks.

Fluffy is a tiger.

Walt is a person who has pets called Huey, Louie and Dewey.

Chunk 5. Formalize in DL, and then in OWL DL

1. Concept definitions

A driver must be adult. A driver is a person who drives vehicles. A lorry driver is a person who drives lorries. A haulage worker is who works for a haulage company or for part of a haulage company. A haulage truck driver is a person who drives trucks and works for part of a haulage company. A van driver is a person who drives vans. A bus driver is a person who drives buses. A white van man is a man who drives white things and vans.

2. Restrictions:

--

3. Properties:

The service number is an integer property with no restricted domain

4. Individuals:

Q123ABC is a van and a white thing.

The42 is a bus whose service number is 42.

Mick is a male who read Daily Mirror and drives Q123ABC.

Chunk 1. Formalisation in DL

$grass \sqsubseteq plant$

$tree \sqsubseteq plant$

$leaf \sqsubseteq \exists partOf . tree$

$dog \sqsubseteq \exists eats . bone$

$sheep \sqsubseteq animal \cap \forall eats . grass$

$giraffe \sqsubseteq animal \cap \forall eats . leaf$

$madCow \equiv cow \cap \exists eats . (brain \cap \exists partOf . sheep)$

$(animal \cup \exists partOf . animal) \cap (plant \cup \exists partOf . plant) \sqsubseteq \perp$

Chunk 2. Formalisation in DL

$bicycle \sqsubseteq vehicle; bus \sqsubseteq vehicle; car \sqsubseteq vehicle; lorry \sqsubseteq vehicle; truck \sqsubseteq vehicle$

$busCompany \sqsubseteq company; haulageCompany \sqsubseteq company$

$elderly \sqsubseteq person \cap adult$

$kid \equiv person \cap young$

$man \equiv person \cap male \cap adult$

$woman \equiv person \cap female \cap adult$

$grownUp \equiv person \cap adult$

$oldLady \equiv person \cap female \cap elderly$

$oldLady \sqsubseteq \exists hasPet. animal \cap \forall hasPet. cat$

$young \cap adult \sqsubseteq \perp$

$hasMother \sqsubseteq hasParent$

$hasFather \sqsubseteq hasParent$

Chunk 3. Formalisation in DL

$magazine \sqsubseteq publication$

$broadsheet \sqsubseteq newspaper$

$tabloid \sqsubseteq newspaper$

$qualityBroadsheet \sqsubseteq broadsheet$

$redTop \sqsubseteq tabloid$

$newspaper \sqsubseteq publication \cap (broadsheet \cup tabloid)$

$whiteVanMan \sqsubseteq \forall reads. tabloid$

$tabloid \cap broadsheet \sqsubseteq \perp$

Chunk 4. Formalisation in DL

$pet \equiv \exists isPetOf.T$

$animal \sqsubseteq \exists eats.T$

$vegetarian \equiv animal \cap \forall eats.\neg animal \cap \forall eats.\neg(\exists partOf.animal)$

$duck \sqsubseteq animal; cat \sqsubseteq animal; tiger \sqsubseteq animal$

$animalLover \equiv person \cap (\geq 3 hasPet)$

$petOwner \equiv person \cap \exists hasPet.animal$

$catLike \equiv person \cap \exists likes.cat; catOwner \equiv person \cap \exists hasPet.cat$

$dogLike \equiv person \cap \exists likes.dog; dogOwner \equiv person \cap \exists hasPet.dog$

$dog \cap cat \sqsubseteq \perp$

Chunk 5. Formalisation in DL

$driver \sqsubseteq adult$

$driver \equiv person \cap \exists drives.vehicle$

$lorryDriver \equiv person \cap \exists drives.lorry$

$haulageWorke \equiv \exists worksFor.(haulageCompany \cup \exists partOf.haulageCompany)$

$haulageTruckDriver \equiv person \cap \exists drives.truck \cap$

$\exists worksFor.(\exists partOf.haulageCompany)$

$vanDriver \equiv person \cap \exists drives.van$

$busDriver \equiv person \cap \exists drives.bus$

$whiteVanMan \equiv man \cap \exists drives.(whiteThing \cap van)$

OWL Example

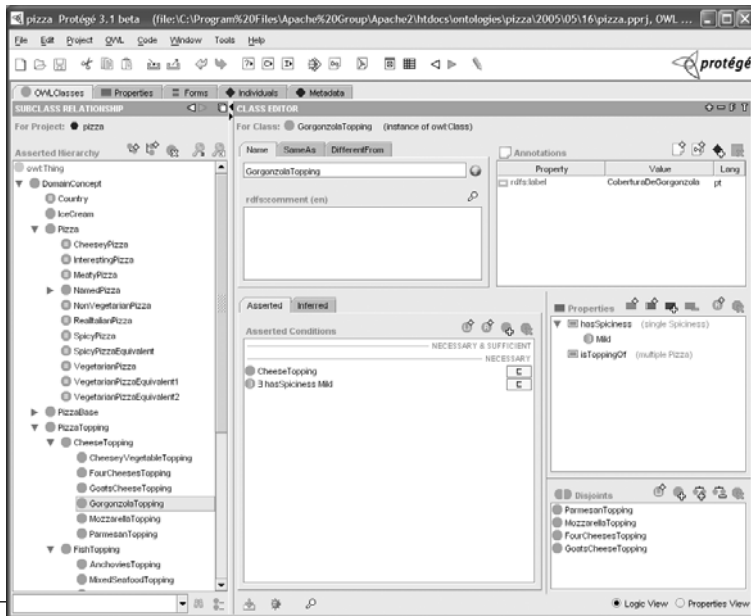


Protégé and the Protégé-OWL plug-in

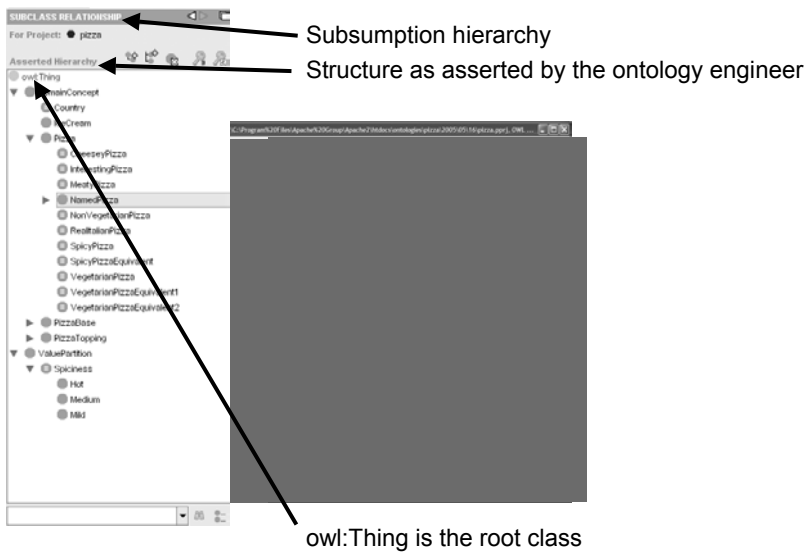
- <http://protege.stanford.edu/>
- **Developed by Stanford Medical Informatics**
- **Features**
 - Software
 - Open source
 - Supports development of plugins to allow backend / interface extensions
 - Large user community (approx 30k)
 - Knowledge representation formalism
 - Core is based on Frames (object oriented) modelling
 - Open architecture that allows other modelling languages to be built on top (e.g., Protégé-OWL plug-in)



Protégé-OWL



Class Hierarchy



Class Editor

Class annotations (for class metadata)

Class name and documentation

Properties "available" to Class

Disjoints widget

Conditions Widget

Class-specific tools (find usage etc)

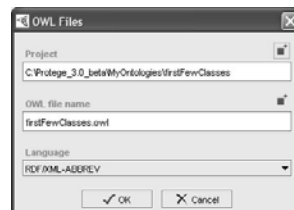


Saving OWL Files

OWL = easy to make mistakes = save regularly



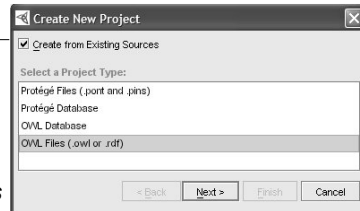
1. *Select File → Save Project As*
A dialog (as shown) will pop up
2. *Select a file directly by clicking the button on the top right*
You will notice that 2 files are created
.prj – the project file
this just stores information about the GUI and the workspace
.owl – the OWL file
this is where your ontology is stored in RDF/OWL format
3. *Select OK*



Loading OWL files

1. *If you only have an OWL file:*
 - **File** → **New Project**
 - Select **OWL Files** as the type
 - Tick **Create from existing sources**
 - **Next** to select the **.owl** file
2. *If you've got a valid project file*:*
 - **File** → **Open Project**
 - select the **.pproj** file

* ie one created on this version of Protégé - the s/w gets updated once every few days, so don't count on it unless you've created it recently- safest to build from the **.owl** file if in doubt



OWL Example



Implement the previous sample ontology with the Protégé-OWL plug-in

- Practice with Protégé and the Protégé-OWL plug-in
- Develop parts of the previous ontology in groups, so that the ontologies can be integrated later

OWL Example



Basic Inference Tasks

- **Subsumption – check knowledge is correct (captures intuitions)**
 - Does C subsume D w.r.t. ontology O ? (in *every* model I of O , $C^I \subseteq D^I$)
- **Equivalence – check knowledge is minimally redundant (no unintended synonyms)**
 - Is C equivalent to D w.r.t. O ? (in *every* model I of O , $C^I = D^I$)
- **Consistency – check knowledge is meaningful (classes can have instances)**
 - Is C satisfiable w.r.t. O ? (there exists *some* model I of O s.t. $C^I \neq \emptyset$)
- **Instantiation and querying**
 - Is x an instance of C w.r.t. O ? (in *every* model I of O , $x^I \in C^I$)
 - Is (x,y) an instance of R w.r.t. O ? (in *every* model I of O , $(x^I, y^I) \in R^I$)
- **All reducible to KB satisfiability or concept satisfiability w.r.t. a KB**
- **Can be decided using highly optimised tableaux reasoners**

Tableaux Algorithms

- **Try to prove satisfiability of a knowledge base**
- **How do they work**
 - They try to build a model of input concept C
 - Tree model property
 - If there is a model, then there is a tree shaped model
 - If no tree model can be found, then input concept unsatisfiable
 - Decompose C syntactically
 - Work on concepts in negation normal form (De Morgan's laws)
 - Use of tableaux expansion rules
 - If non-deterministic rules are applied, then there is search
 - Stop (and backtrack) if clash
 - E.g. $A(x), \neg A(x)$
 - Blocking (cycle check) ensures termination for more expressive logics
- **The algorithm finishes when no more rules can be applied or a conflict is detected**

Tableaux rules for ALC and for transitive roles

$x \bullet \{C_1 \sqcap C_2, \dots\}$	$\rightarrow \sqcap$	$x \bullet \{C_1 \sqcap C_2, C_1, C_2, \dots\}$
$x \bullet \{C_1 \sqcup C_2, \dots\}$	$\rightarrow \sqcup$	$x \bullet \{C_1 \sqcup C_2, C, \dots\}$ for $C \in \{C_1, C_2\}$
$x \bullet \{\exists R.C, \dots\}$	$\rightarrow \exists$	$x \bullet \{\exists R.C, \dots\}$ R \downarrow $y \bullet \{C\}$
$x \bullet \{\forall R.C, \dots\}$ R \downarrow $y \bullet \{\dots\}$	$\rightarrow \forall$	$x \bullet \{\forall R.C, \dots\}$ R \downarrow $y \bullet \{C, \dots\}$
$x \bullet \{\forall R.C, \dots\}$ R \downarrow $y \bullet \{\dots\}$	$\rightarrow \forall_+$	$x \bullet \{\forall R.C, \dots\}$ R \downarrow $y \bullet \{\forall R.C, \dots\}$

Tableaux examples and exercises

- **Example**

- $\exists R.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)$

- **Exercise 1**

- $\exists R.(\exists R.D) \wedge \exists S.\neg D \wedge \forall S.(\exists R.D)$

- **Exercise 2**

- $\exists R.(C \vee D) \wedge \forall R.\neg C \wedge \neg \exists R.D$

OWL Example



Develop a sample ontology in the domain of people, pets, vehicles, and newspapers

- Understand the basic reasoning mechanisms of OWL DL

Subsumption

Automatic classification: an ontology built collaboratively

Instance classification

Detecting redundancy

Consistency checking: unsatisfiable restrictions in a Tbox (are the classes coherent?)

Tableaux Algorithm — Example

Test satisfiability of $\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)$ where R is a **transitive** role

Tableaux Algorithm — Example

Test satisfiability of $\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)$ where R is a **transitive** role

$$\mathcal{L}(w) = \{\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)\}$$

w

Tableaux Algorithm — Example

Test satisfiability of $\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)$ where R is a **transitive** role

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w

Tableaux Algorithm — Example

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$$\mathcal{L}(w) = \{\exists S.C, \forall S.(\neg C \sqcup \neg D), \exists R.C, \forall R.(\exists R.C)\}$$

w

Tableaux Algorithm — Example

Test satisfiability of $\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)$ where R is a **transitive** role

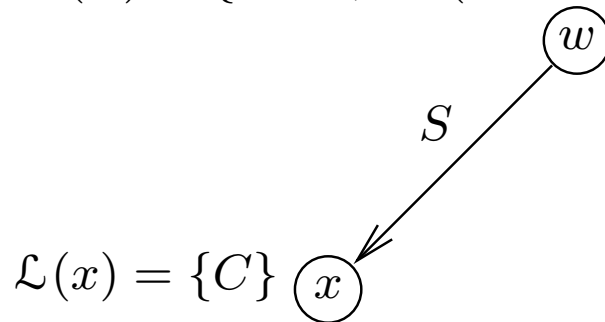
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Tableaux Algorithm — Example

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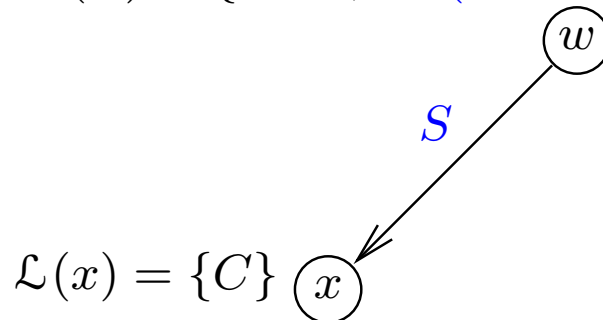
$$\mathcal{L}(w) = \{\exists S.C, \forall S.(\neg C \sqcup \neg D), \exists R.C, \forall R.(\exists R.C)\}$$



Tableaux Algorithm — Example

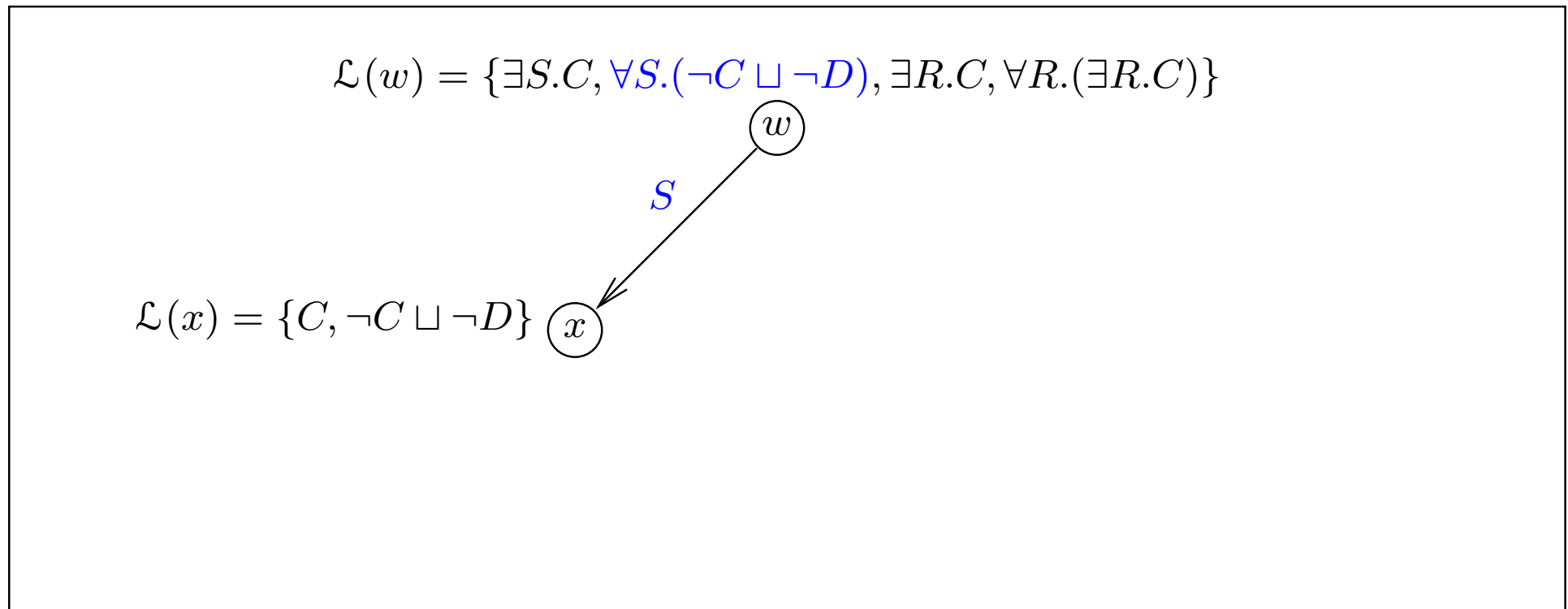
Test satisfiability of $\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)$ where R is a **transitive** role

$$\mathcal{L}(w) = \{\exists S.C, \forall S.(\neg C \sqcup \neg D), \exists R.C, \forall R.(\exists R.C)\}$$



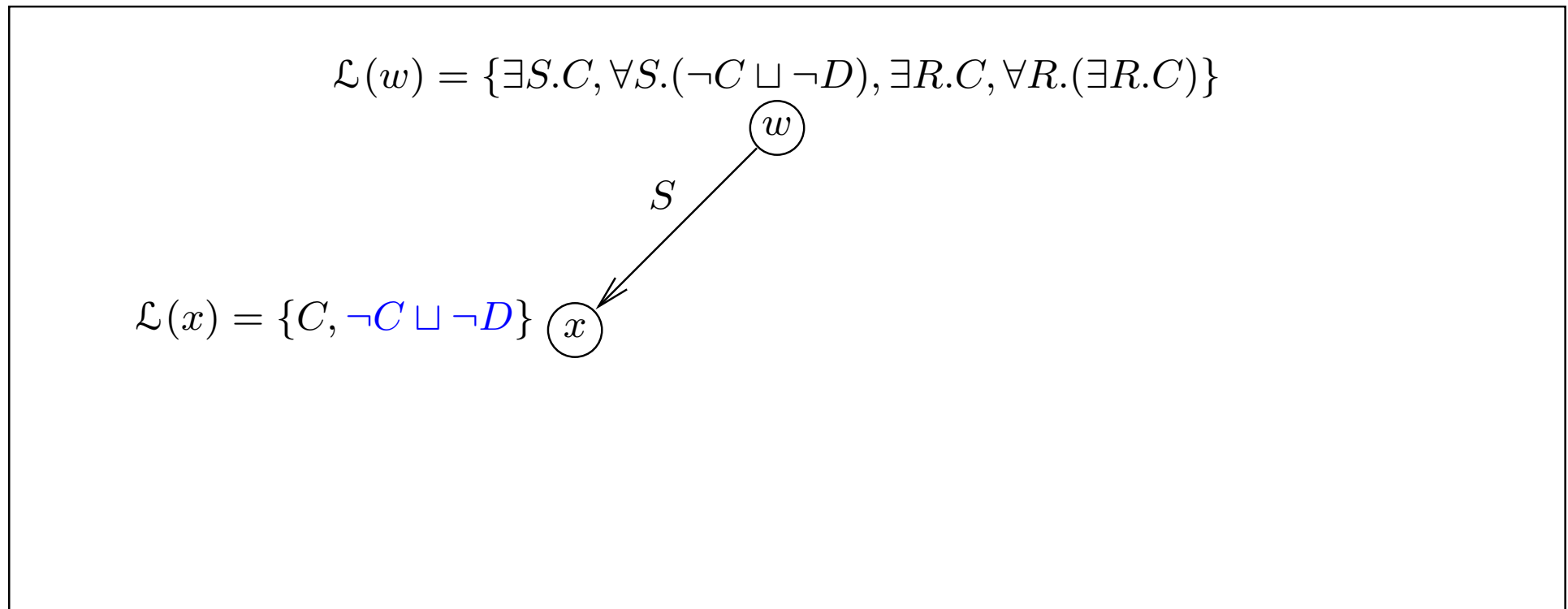
Tableaux Algorithm — Example

Test satisfiability of $\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)$ where R is a **transitive** role



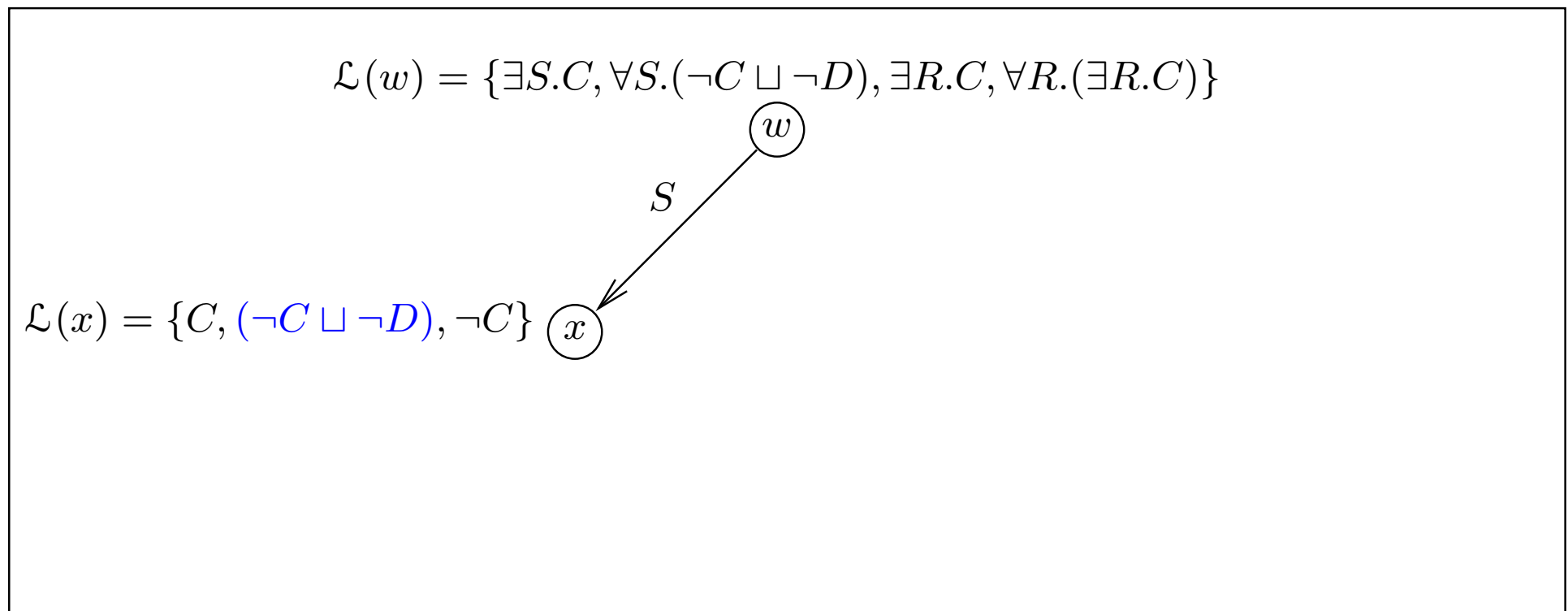
Tableaux Algorithm — Example

Test satisfiability of $\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)$ where R is a **transitive** role



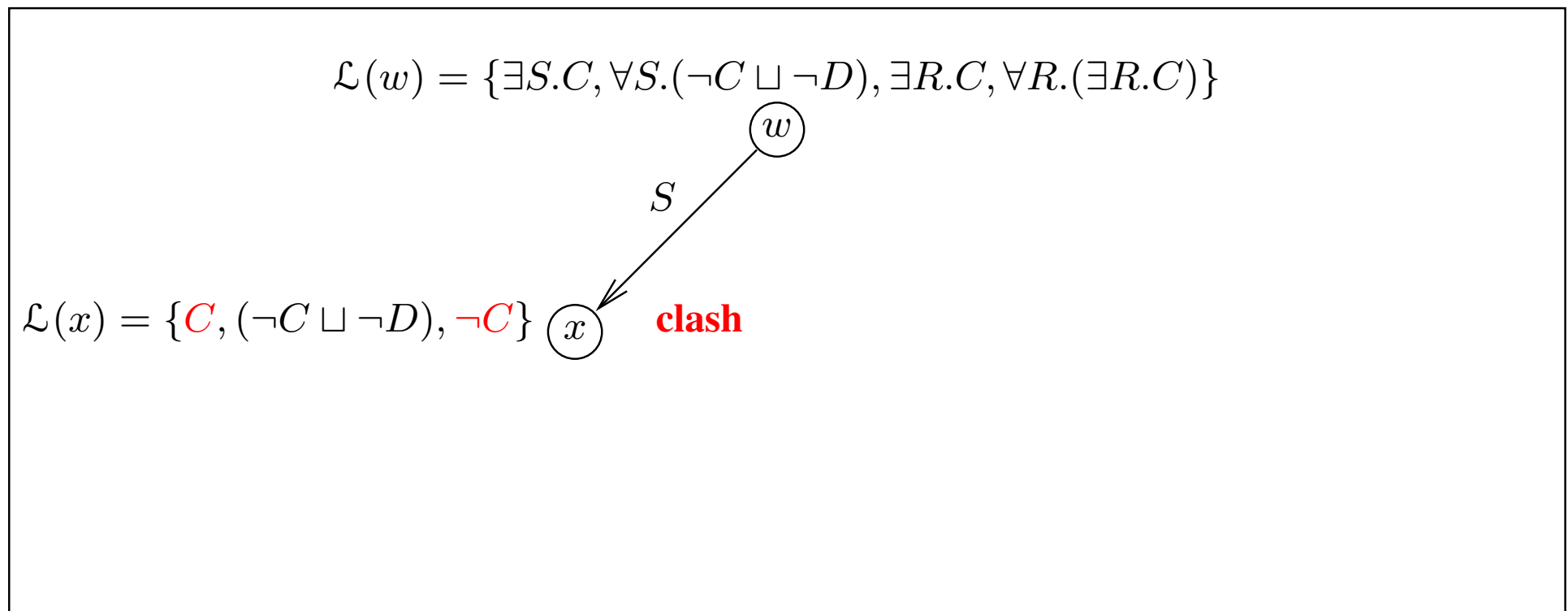
Tableaux Algorithm — Example

Test satisfiability of $\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)$ where R is a **transitive** role



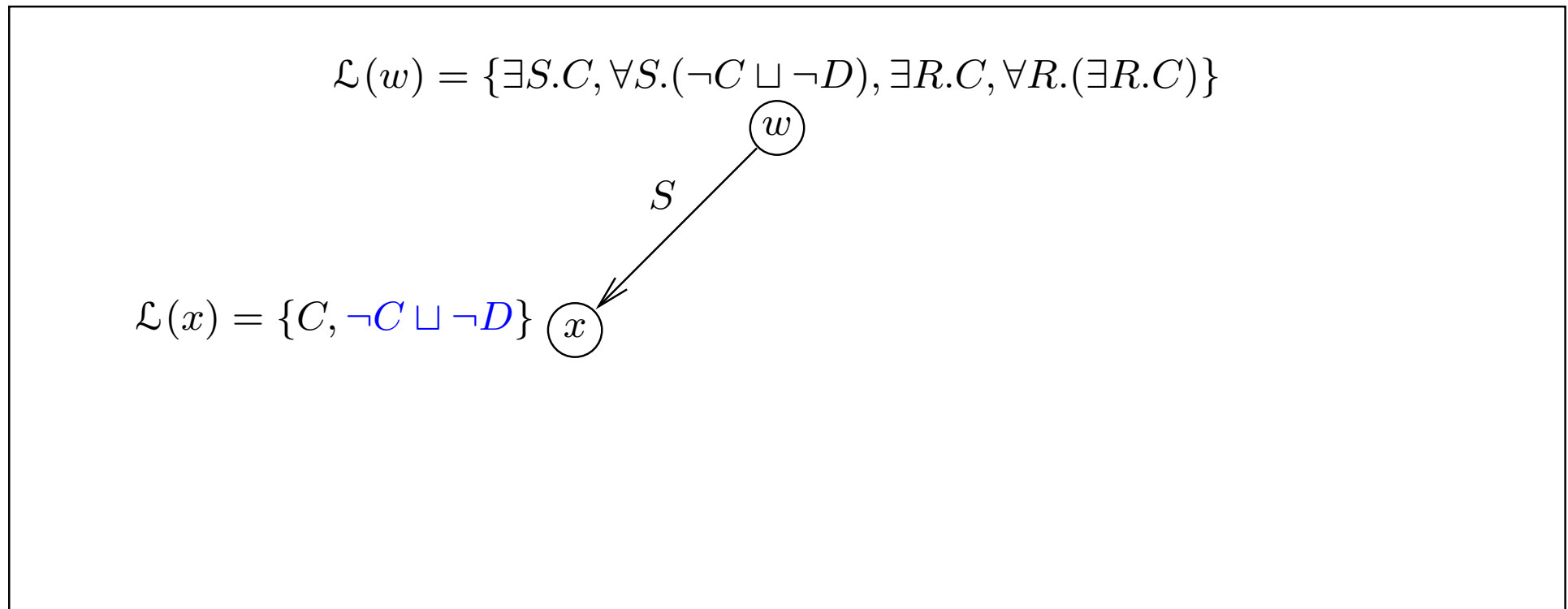
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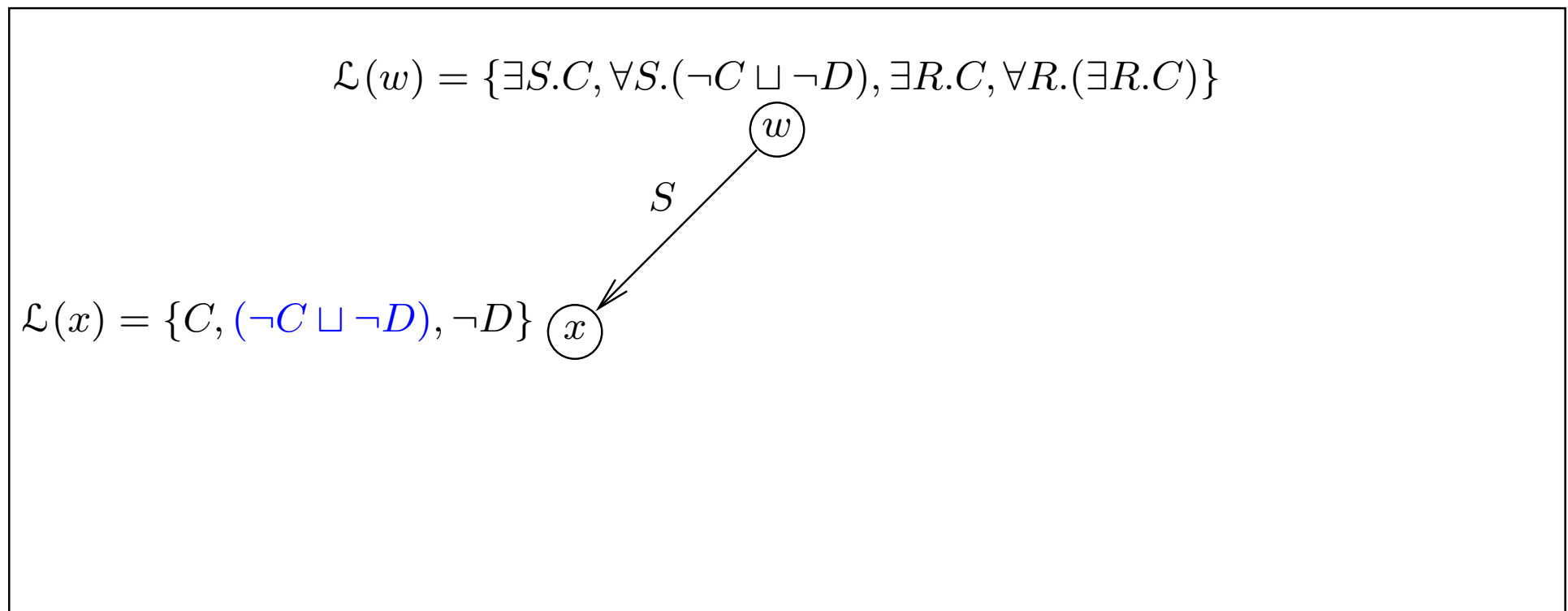
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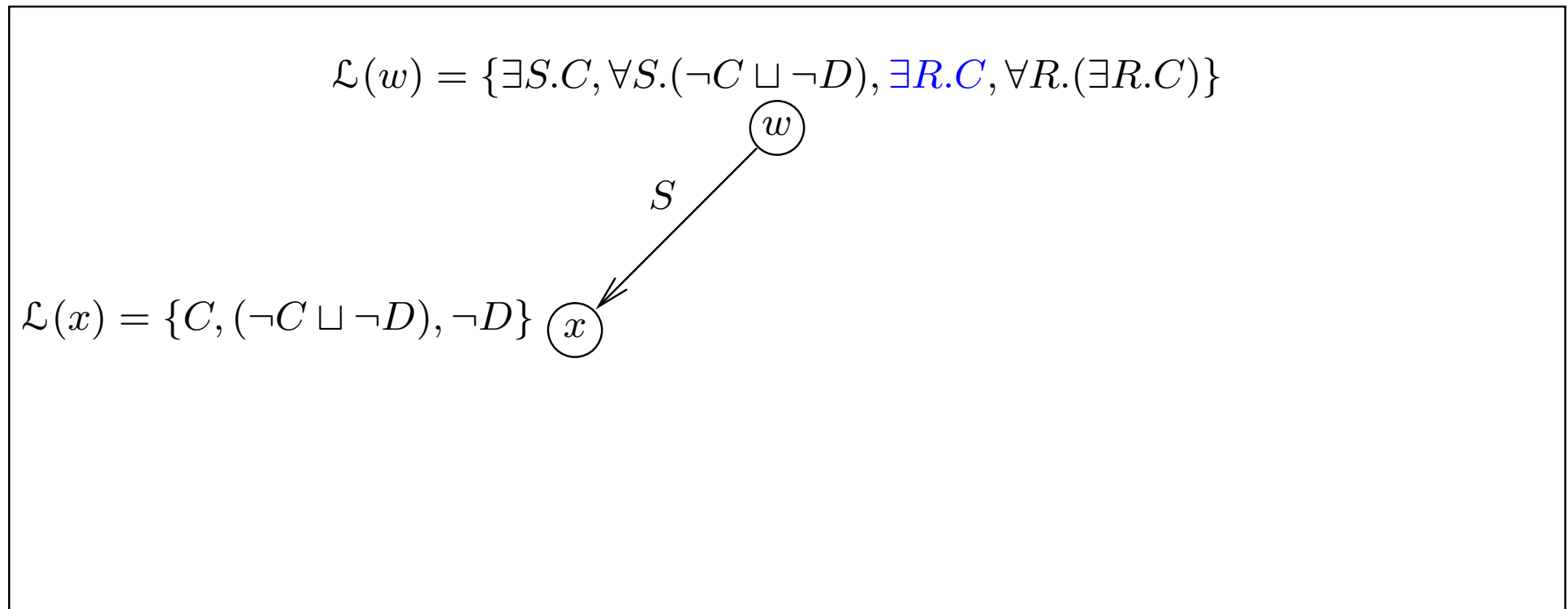
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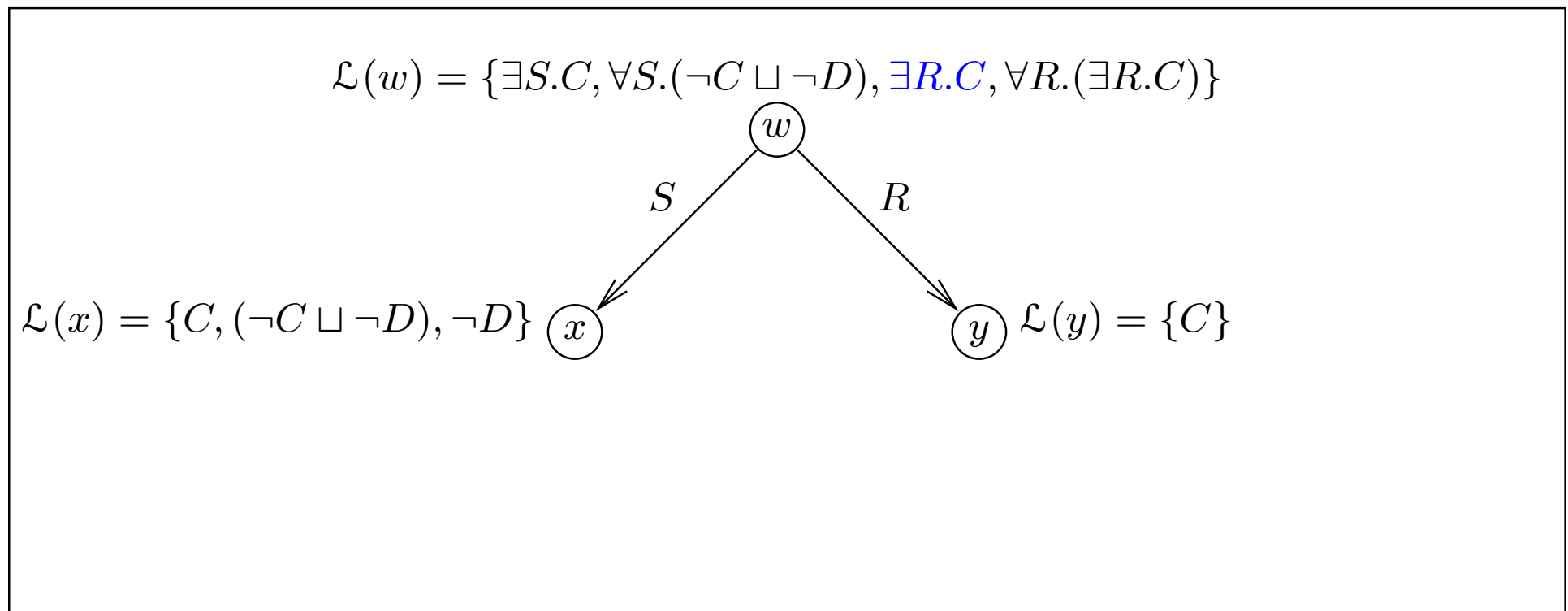
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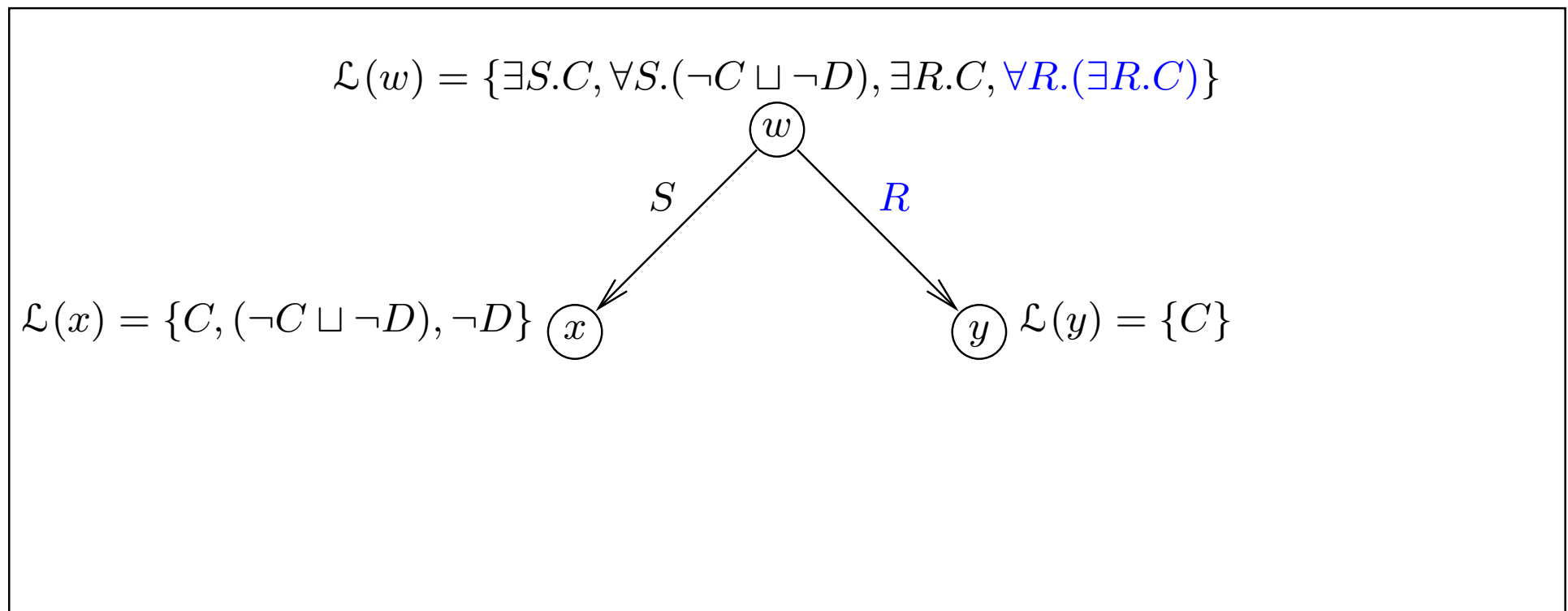
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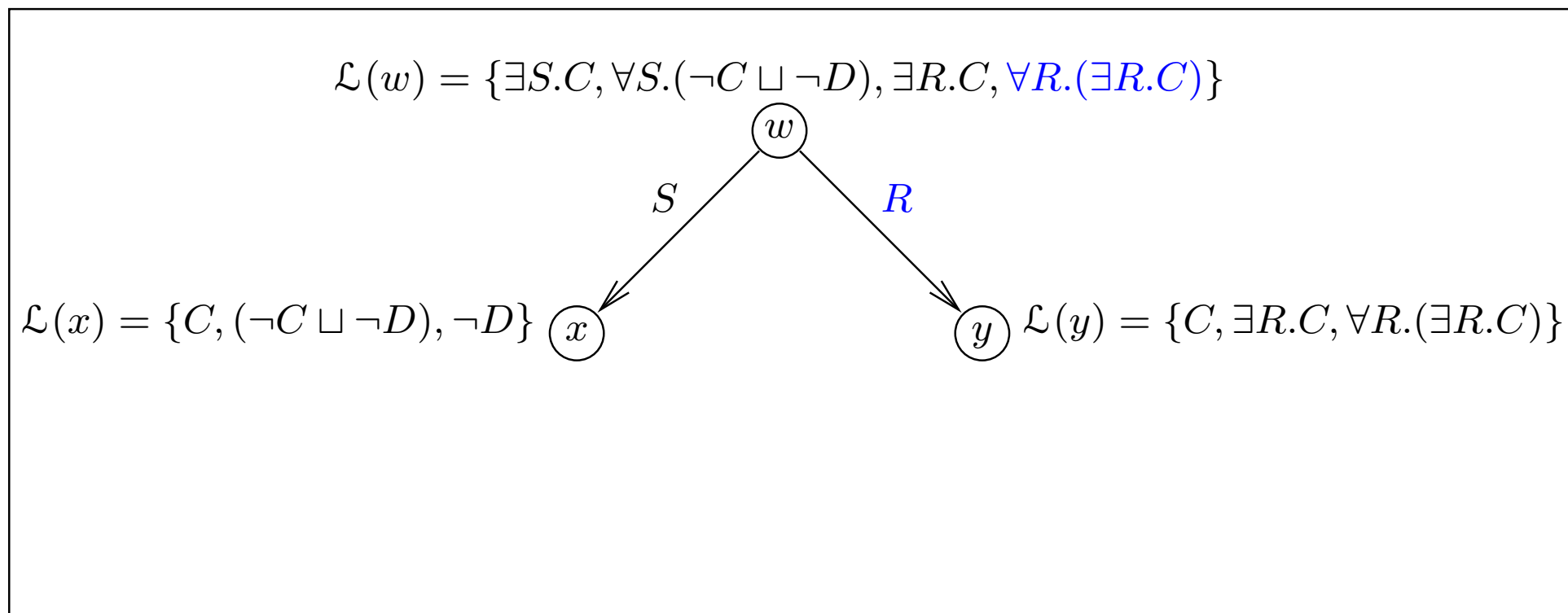
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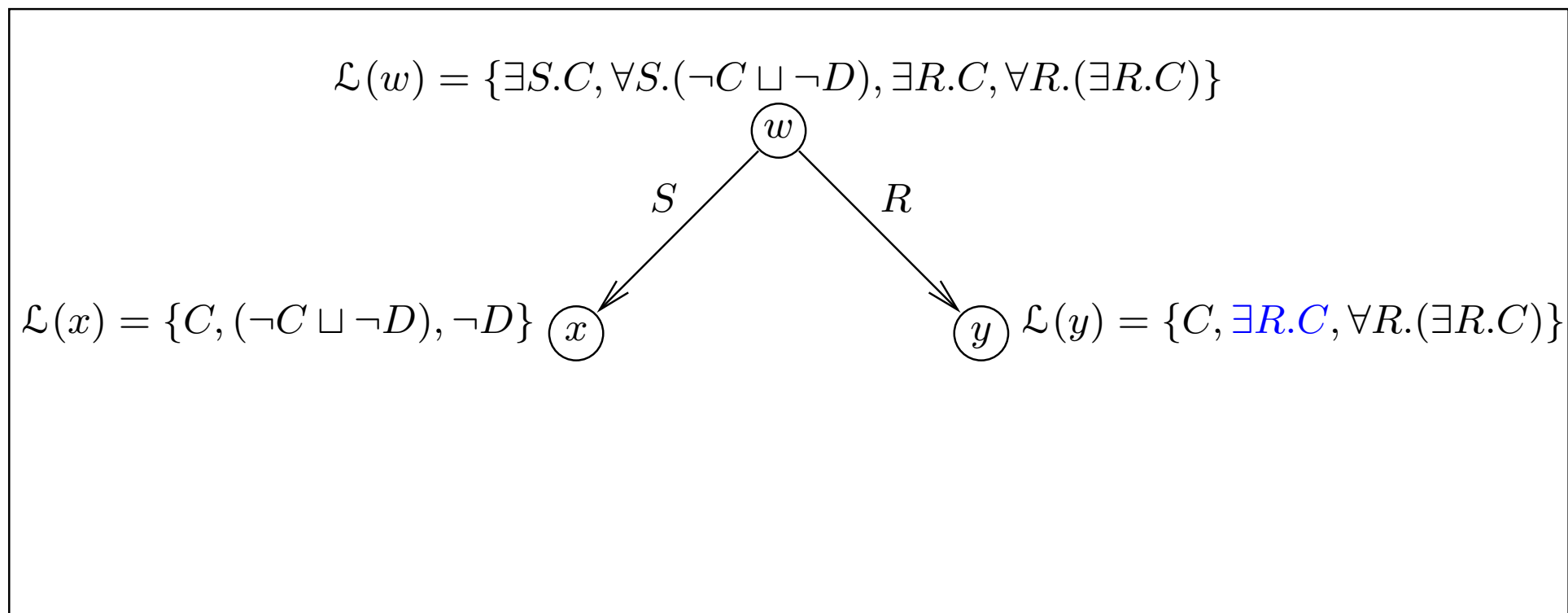
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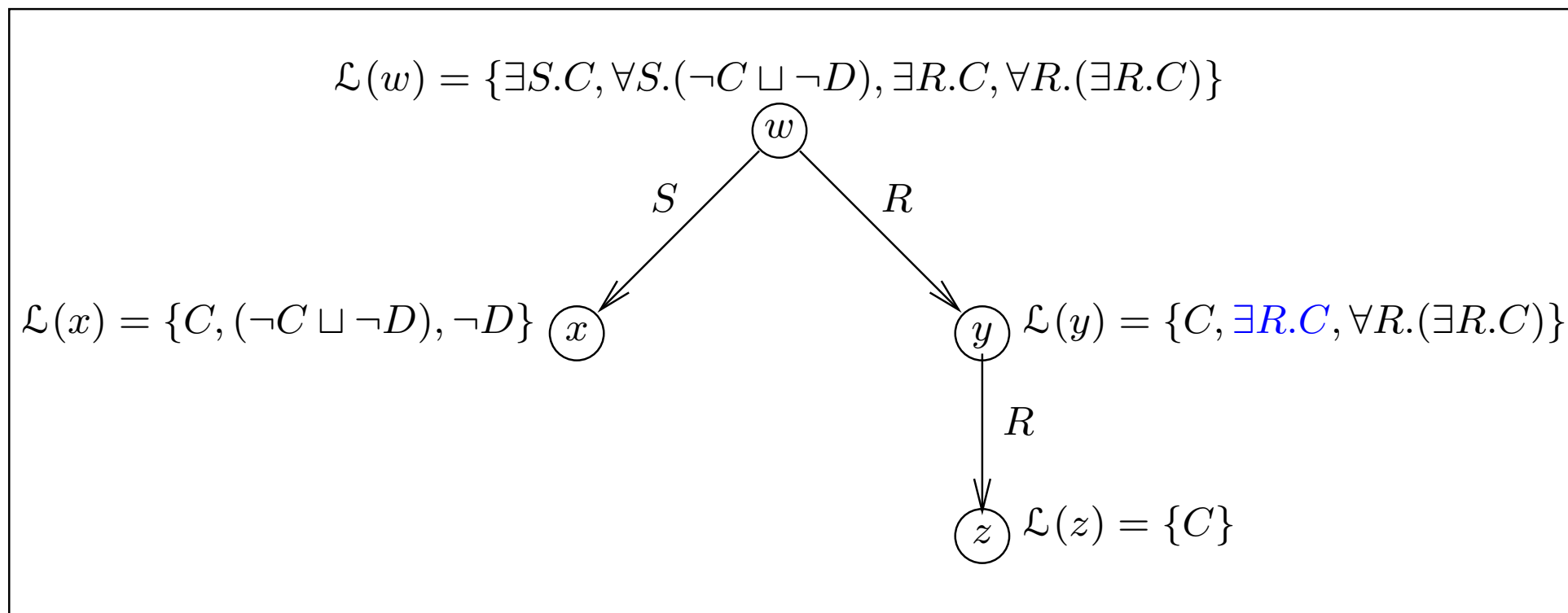
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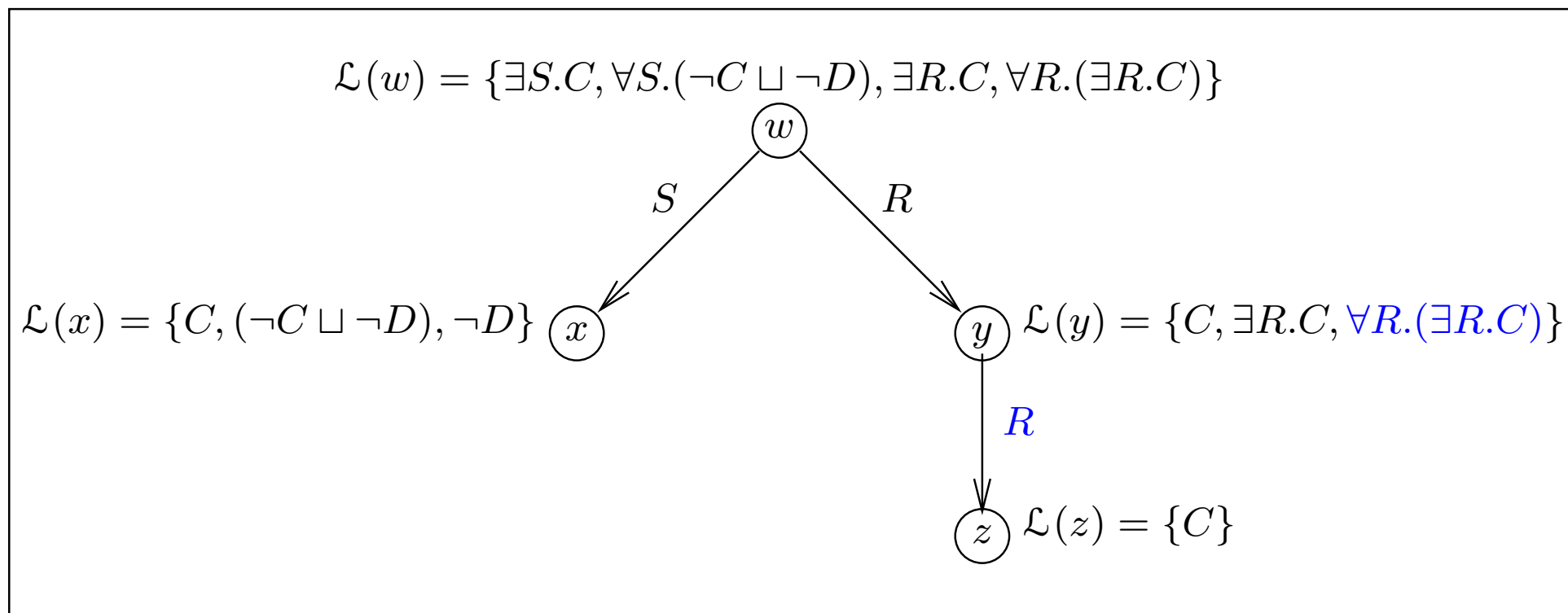
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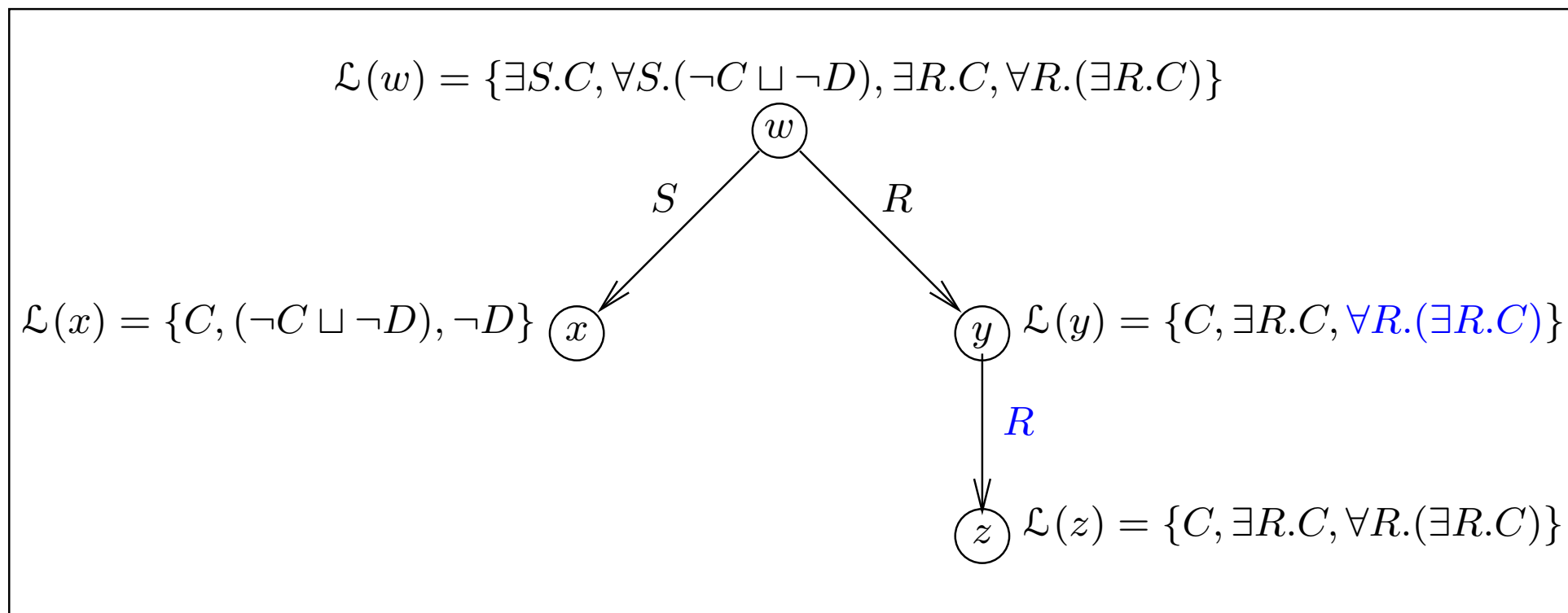
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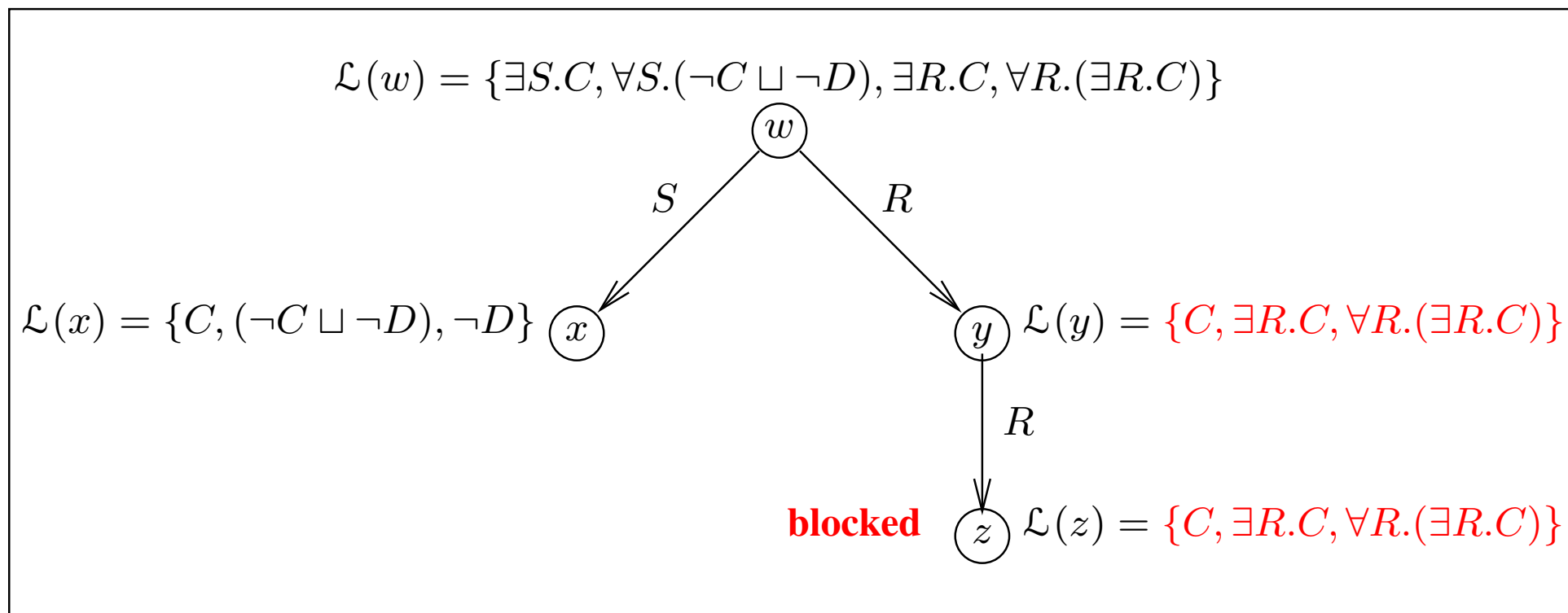
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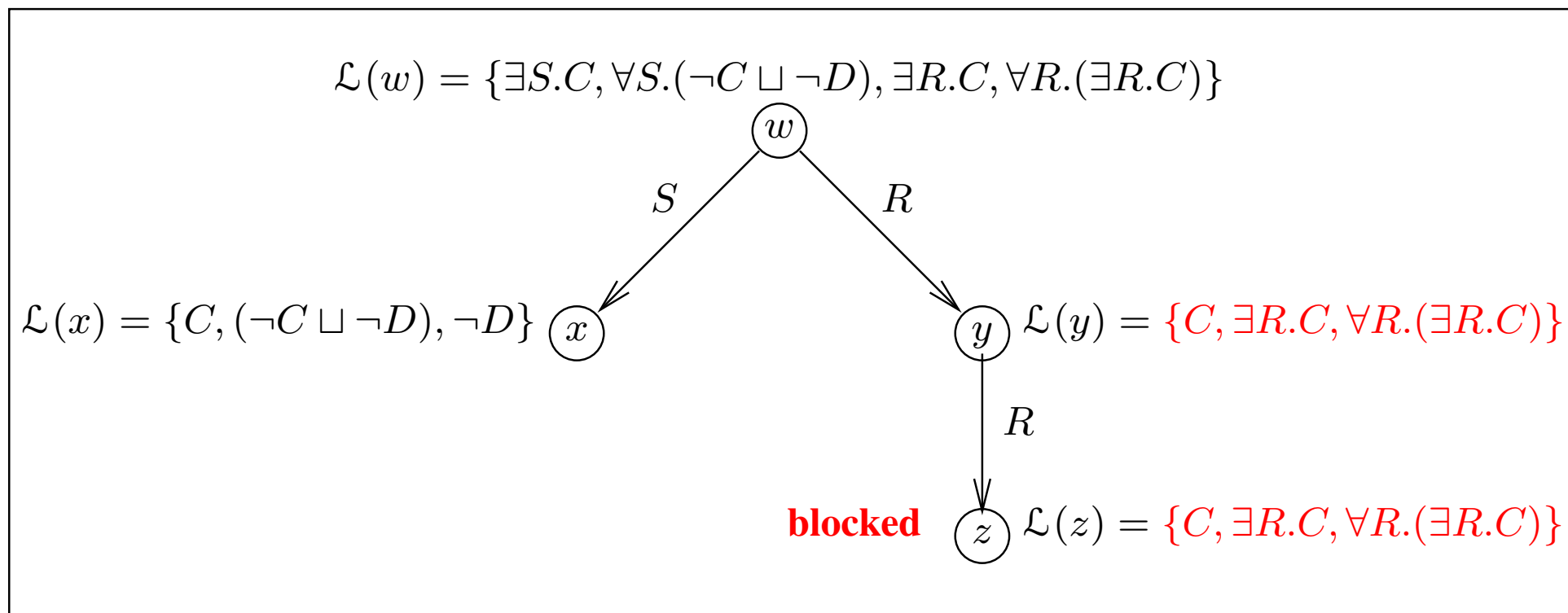
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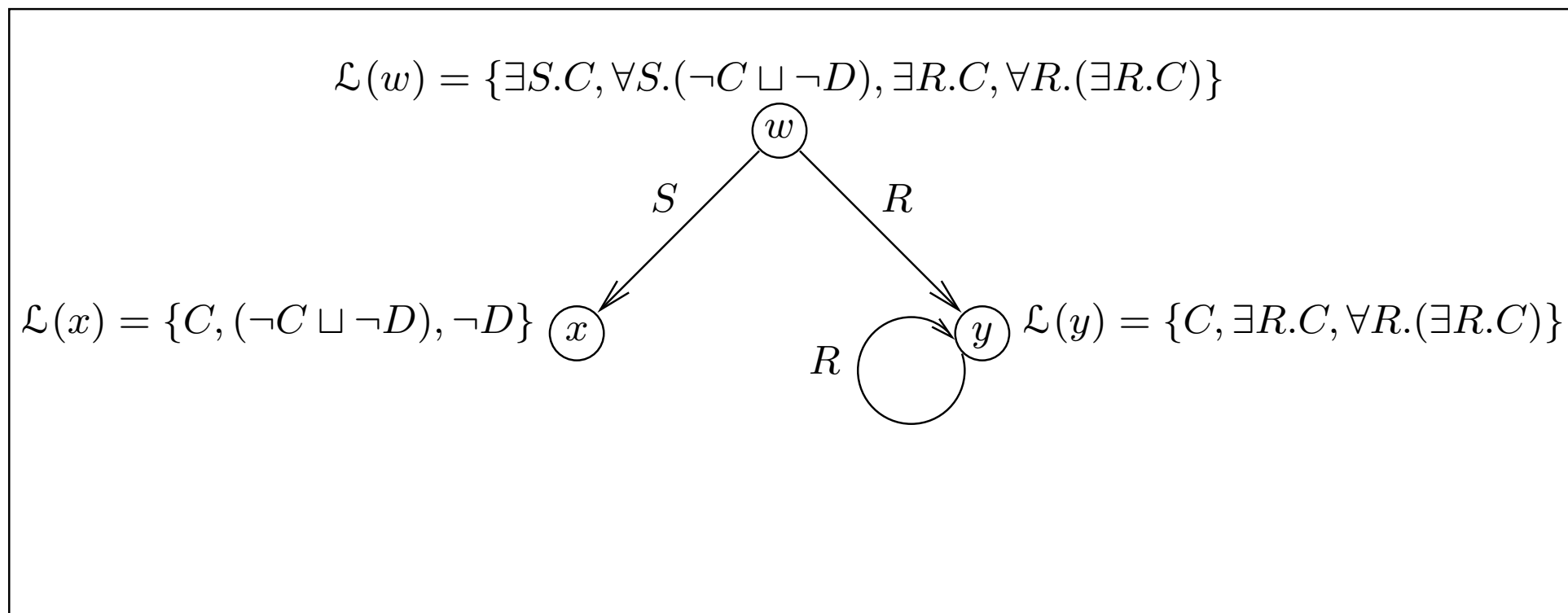
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Concept is **satisfiable**: \mathbb{T} corresponds to **model**

Tableaux Algorithm — Example

Test satisfiability of $\exists S.C \sqcap \forall S.(\neg C \sqcup \neg D) \sqcap \exists R.C \sqcap \forall R.(\exists R.C)$ where R is a **transitive** role

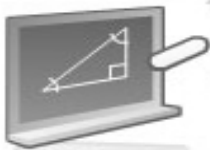


Concept is **satisfiable**: \mathbb{T} corresponds to **model**

Tableaux examples and exercises

- **Example**
 - $\exists S.C \cap \forall S.(\neg C \sqcup \neg D) \cap \exists R.C \cap \forall R.(\exists R.C)$
- **Exercise 1**
 - $\exists R.(\exists R.D) \wedge \exists S.\neg D \wedge \forall S.(\exists R.D)$
- **Exercise 2**
 - $\exists R.(C \vee D) \wedge \forall R.\neg C \wedge \neg \exists R.D$

OWL Example



Develop a sample ontology in the domain of people, pets, vehicles, and newspapers

- Understand the basic reasoning mechanisms of OWL DL

Subsumption

Automatic classification: an ontology built collaboratively

Instance classification

Detecting redundancy

Consistency checking: unsatisfiable restrictions in a Tbox (are the classes coherent?)

Interesting results (I). Automatic classification

And old lady is a person who is elderly and female.

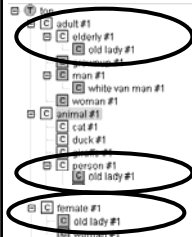
Old ladies must have some animal as pets and all their pets are cats.

$$\text{elderly} \subseteq \text{person} \cap \text{adult}$$

$$\text{woman} \equiv \text{person} \cap \text{female} \cap \text{adult}$$

$$\text{catOwner} \equiv \text{person} \cap \exists \text{hasPet.cat}$$

$$\text{oldLady} \equiv \text{person} \cap \text{female} \cap \text{elderly}$$

$$\text{oldLady} \subseteq \exists \text{hasPet.animal} \cap \forall \text{hasPet.cat}$$


We obtain:

Old ladies must be women.

Every old lady must have a pet cat

Hence, every old lady must be a cat owner

$$\text{oldLady} \subseteq \text{woman} \cap \text{elderly} \cap \text{catOwner}$$


Interesting results (II). Instance classification

A pet owner is a person who has animal pets

Old ladies must have some animal as pets and all their pets are cats.

Has pet has domain person and range animal

Minnie is a female, elderly, who has a pet called Tom.

$$\text{petOwner} \equiv \text{person} \cap \exists \text{hasPet.animal}$$

$$\text{oldLady} \subseteq \exists \text{hasPet.animal} \cap \forall \text{hasPet.cat}$$

$$\text{hasPet} \subseteq (\text{person}, \text{animal})$$

$$\text{Minnie} \in \text{female} \cap \text{elderly}$$

$$\text{hasPet}(\text{Minnie}, \text{Tom})$$

We obtain:

Minnie is a person

Hence, Minnie is an old lady

Hence, Tom is a cat

$$\text{Minnie} \in \text{person}; \text{Tom} \in \text{animal}$$

$$\text{Minnie} \in \text{petOwner}$$

$$\text{Minnie} \in \text{oldLady}$$

$$\text{Tom} \in \text{cat}$$

Interesting results (III). Instance classification and redundancy detection

An animal lover is a person who has at least three pets

Walt is a person who has pets called Huey, Louie and Dewey.

$animalLover \equiv person \cap (\geq 3hasPet)$

$Walt \in person$

$hasPet(Walt, Huey)$

$hasPet(Walt, Louie)$

$hasPet(Walt, Dewey)$

We obtain:

Walt is an animal lover

Walt is a person is redundant

$Walt \in animalLover$

Interesting results (IV). Instance classification

A van is a type of vehicle

A driver must be adult

A driver is a person who drives vehicles

A white van man is a man who drives vans and white things

White van mans must read only tabloids

Q123ABC is a white thing and a van

Mick is a male who reads Daily Mirror and drives Q123ABC

$van \sqsubseteq vehicle$

$driver \sqsubseteq adult$

$driver \equiv person \cap \exists drives. vehicle$

$whiteVanMan \equiv man \cap \exists drives. (van \cap whiteThing)$

$whiteVanMan \sqsubseteq \forall reads. tabloid$

$Q123ABC \in whiteThing \cap van$

$Mick \in male$

$reads(Mick, DailyMirror)$

$drives(Mick, Q123ABC)$

We obtain:

Mick is an adult

Mick is a white van man

Daily Mirror is a tabloid

$Mick \in adult$

$Mick \in whiteVanMan$

$DailyMirror \in tabloid$

Interesting results (V). Consistency checking

Cows are vegetarian.

A vegetarian is an animal that does not eat animals nor parts of animals.

A mad cow is a cow that eats brains that can be part of a sheep

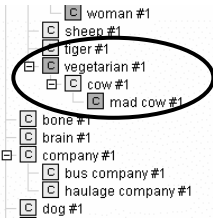
$cow \sqsubseteq vegetarian$

$vegetarian \equiv animal \cap \forall eats. \neg animal \cap$

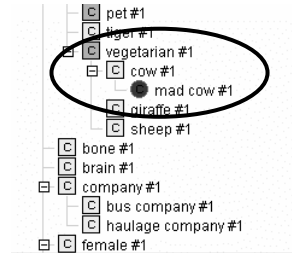
$\forall eats. \neg (\exists partOf. animal))$

$madCow \equiv cow \cap \exists eats. (brain \cup \exists partOf. sheep)$

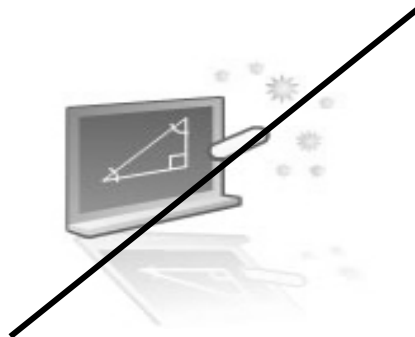
$(animal \cup \exists partOf. animal) \cap (plant \cup \exists partOf. plant) \sqsubseteq \perp$



We obtain:
Mad cow is unsatisfiable



OWL Example



When to use a classifier

1. At author time (pre-coordination): As a compiler

- Ontologies will be delivered as “pre-coordinated” ontologies to be used without a reasoner
- To make extensions and additions quick, easy, and responsive, distribute developments, empower users to make changes
- Part of an ontology life cycle

2. At delivery time (post-coordination): as a normalisation service

- Many fixed ontologies are too big and too small
 - Too big to find things; too small to contain what you need
 - Create them on the fly
- Part of an ontology service

3. At application time (inference): as a reasoner

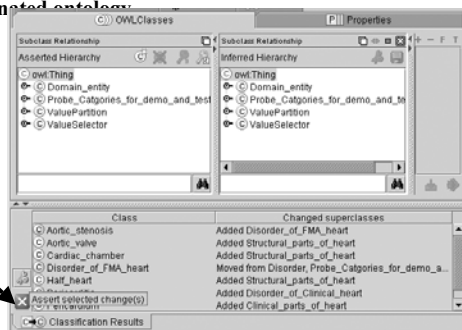
- Decision support, query optimisation, schema integration, etc.
- Part of a reasoning service



1. Pre-coordinated delivery: classifier as compiler

- **Develop an ontology**
 - A classifier can be used to detect and correct inconsistencies
- **Classify the ontology**
- **Commit classifier results to a pre-coordinated ontology:**

Assert (“Commit”) changes inferred by classifier

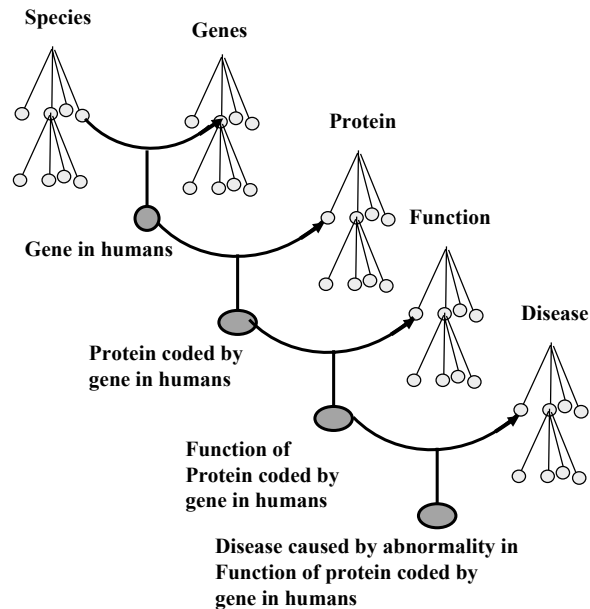
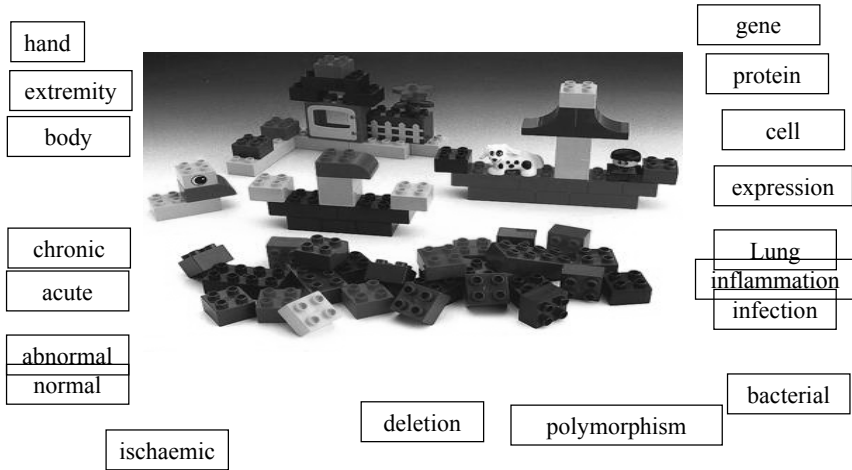


- **Deliver it**
 - In OWL-Lite or RDFS
- **Use RDQL, SPARQL, or your favourite RDF(S) query tool**



2. Post Coordination: classifier as a service

- Logic based ontologies act as a conceptual lego
 - Modularisation/Normalisation is needed to make them easier to maintain



Logical Constructs build complex concepts from modularised primitives

Disease caused by abnormality in Function of protein coded by gene in humans

Rationale for Normalisation/Modularisation

- **Normalisation in order to make large ontologies easier to maintain**
 - Trees (classes have only one parent) are easier to maintain
 - Directed Acyclic Graphs (classes can have several parents) can be derived with a reasoner from the trees
- **Objective: derive explicit distinctions between modules**
 - Primitives are opaque to the reasoner
 - Information implicit in primitive names cannot contribute to modularisation
 - Primitives are indivisible to both human and reasoner
 - Each primitive should represent a single notion
 - Therefore, each primitive must belong to exactly one module
 - If a primitive belongs to two modules, they are not modular.
 - If a primitive belongs to two modules, it probably conflates two notions
 - Therefore concentrate on the “primitive skeleton” of the domain ontology

Normalisation Criteria (I)

1. The skeleton should consist of disjoint trees

- Every primitive concept should have exactly one primitive parent
 - All multiple hierarchies are the result of inference by reasoner

2. No hidden changes in meaning

- Each branch should be homogeneous and logical (“Aristotelian”)
 - Hierarchical principle should be subsumption
 - Otherwise we are “lying to the logic”
 - The criteria for differentiation should follow consistent principles in each branch
- Example of non-homogeneous taxonomy (from The Celestial Emporium of Benevolent Knowledge, Borges)
 - “On those remote pages it is written that animals are divided into:
 - a. those that belong to the Emperor
 - b. embalmed ones
 - c. those that are trained
 - d. suckling pigs
 - e. mermaids
 - f. fabulous ones
 - g. stray dogs
 - h. those that are included in this classification
 - i. those that tremble as if they were mad
 - j. innumerable ones
 - k. those drawn with a very fine camel's hair brush
 - l. others
 - m. those that have just broken a flower vase
 - n. those that resemble flies from a distance”

Normalisation Criteria (II)

3. Distinguish “Self-standing” and “Refining” Concepts

- Self-standing concepts (person, idea, plant, committee, belief, etc.)
 - Roughly Welty & Guarino’s “sortals”
- Refining concepts – depend on self-standing concepts (mild|moderate|severe, hot|cold, etc.)
 - Roughly Welty & Guarino’s non-sortals, Smith’s “fiat partitions”, Value Types for engineers

3a. Self-standing primitives should be globally disjoint & open

- Primitives are disjoint
 - If primitives overlap, the overlap conceals implicit information
- A list of self-standing primitives can never be guaranteed complete
 - How many kinds of person? of plant? of committee? of belief?
 - Can’t infer: Parent & ¬sub1 &...& ¬subn-1 → subn

3b. Refining primitives should be locally disjoint & closed

- Individual values must be disjoint, but can be hierarchical
 - e.g. “very hot”, “moderately severe”
- Each list can be guaranteed to be complete
- “Value partitions” themselves need not be disjoint
 - “being hot” is not disjoint from “being severe”
 - Allowing Valuetypes to overlap is a useful trick, e.g. restriction has_state someValuesFrom (severe and hot)

Normalisation Criteria (III)

4. Axioms

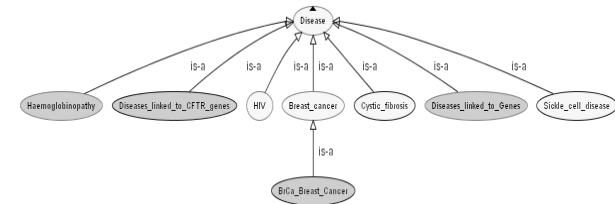
- No axiom should denormalise the ontology
- No axiom should imply that a primitive is part of more than one branch of primitive skeleton
 - If all primitives are disjoint, any such axioms will make that primitive unsatisfiable
 - A partial test for normalisation:
 - Create random conjunctions of primitives which do not subsume each other.
 - If any are satisfiable, the ontology is not normalised

Consequences

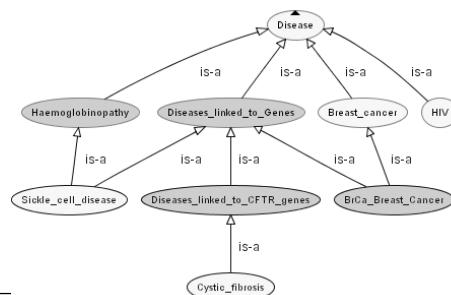
- All self-standing primitives are disjoint
- All multiple classification is inferred
- For any two primitive self-standing classes, either one subsumes the other or they are disjoint
- Every self standing concept is part of exactly one primitive branch of the skeleton
- Every self-standing concept has exactly one most specific primitive ancestor
- Primitives introduced by a conjunction of one class and a boolean combination of zero or more restrictions
 - Tree subclass-of Plant and restriction isMadeOf someValuesFrom Wood
 - Resort subclass-of Accommodation restriction isIntendedFor someValueFrom Holidays
- Use of axioms limited (outside of skeleton construction). The following are a safe but not exhaustive guide:
 - The right side of subclass axioms limited to restrictions
 - Both sides of disjointness axioms limited to restrictions
 - No equivalence axioms with primitives on either side

Example (I)

- Build a simple tree

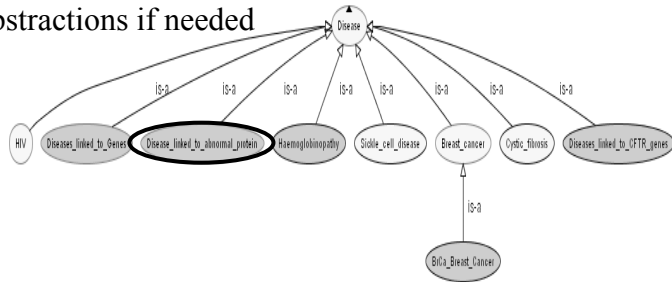


- Let the classifier organise it and check consistency

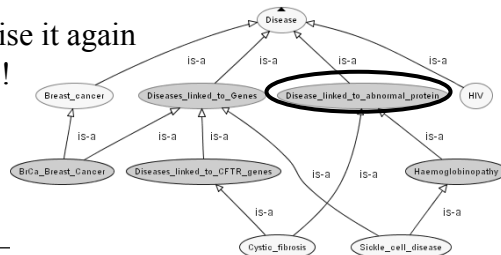


Example (II)

- Add more abstractions if needed



- Let the classifier organise it again and check consistency!!



3. Inference

- We have already seen some examples
 - And we will see some more in session 4

Exercise



•Objective

- Apply design principles to normalise/modularise the OWL ontology developed in the previous exercises

•Tasks

- Identify disjoint trees in the ontology developed.
- Distinguish self-standing and refining classes.
- Identify axioms.
- Go back to the first task if the ontology is not normalised yet.

OWL Classifier limitations

- **Numbers and strings**
 - Simple concrete data types in spec
 - User defined XML data types enmeshed in standards disputes
 - No standard classifier deals with numeric ranges
 - Although several experimental ones do
- **is-part-of and has-part**
 - Totally doubly-linked structures scale horridly
- **Handling of individuals**
 - Variable with different classifiers
 - oneOf works badly with all classifiers at the moment



How can we implement ontologies? Ontology languages

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